


```

CCCCCCCCC DDDDDDDDD DDDDDDDDD LL IIIIII 88888888
CCCCCCCCC DDDDDDDDD DDDDDDDDD LL IIIIII 88888888
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88888888
CC          DD          DD          LL          88888888
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CC          DD          DD          LL          88          88
CCCCCCCCC DDDDDDDDD DDDDDDDDD LLLLLLLLLL IIIIII 88888888
CCCCCCCCC DDDDDDDDD DDDDDDDDD LLLLLLLLLL IIIIII 88888888

```

The image displays three 10x10 grids illustrating the formation of the number 7 using different digits:

- Grid 1 (Left):** The number 7 is formed using the digit 8. The top row is all 8s. The left column is all 8s. The bottom row is all 8s. The right column is all 8s. The middle row is all 8s.
- Grid 2 (Middle):** The number 7 is formed using the digit 3. The top row is all 3s. The left column is all 3s. The bottom row is all 3s. The right column is all 3s. The middle row is all 3s.
- Grid 3 (Right):** The number 7 is formed using the digit 2. The top row is all 2s. The left column is all 2s. The bottom row is all 2s. The right column is all 2s. The middle row is all 2s.

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++
TITLE: CDDLIB

CDD Bliss Library

VERSION: 'V04-000'

FACILITY: Common Data Dictionary

ABSTRACT:

This module is the library file used for compiling all other
modules in the CDD facility.

ENVIRONMENT:

AUTHOR: Jeff East, 22-Jan-80

MODIFIED BY:

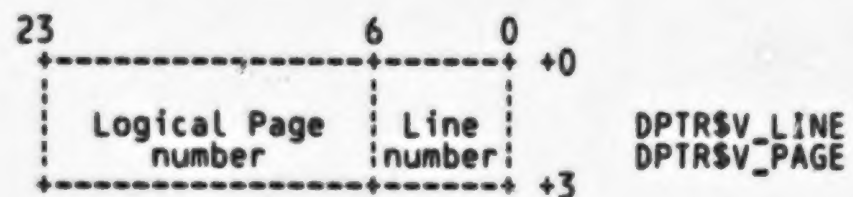
7-May-81 (JAE) Added \$IO_SYNC macro.

--
XTITLE 'CDD Bliss Library'

```
%SBTTL      'STRUCTURE DEFINITIONS'
```

```
STRUCTURE DEFINITIONS
```

```
On-disk Pointer Structure
```

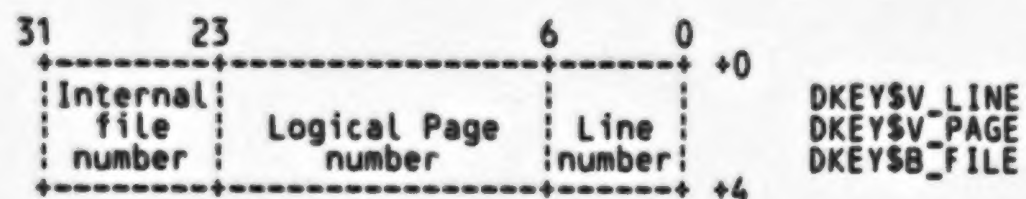


```
LITERAL  
DPTRSS_BLOCK_LENGTH = 3;
```

```
MACRO  
$DPTR = BLOCK[DPTRSS_BLOCK_LENGTH,BYTE] FIELD (DPTRSZ_FIELDS)  
%;
```

```
FIELD DPTRSZ_FIELDS =  
SET  
    DPTRSV_VALUE      = [0, 0, 24, 0],  
    DPTRSV_LINE       = [0, 0, 7, 0],  
    DPTRSV_PAGE       = [0, 7, 17, 0]  
TES;
```


In-core Disk Pointer Structure



When a disk pointer is being passed around routines, it is passed as a Disk Key (DKEY), rather than just a disk pointer.

Disk Keys include all the information of a disk pointer, but also include a pointer to their file's FCB.

LITERAL

DKEYSS_BLOCK_LENGTH = 4;

MACRO

SDKEY = BLOCK[DKEYSS_BLOCK_LENGTH, BYTE] FIELD (DKEYSZ_FIELDS)
%;

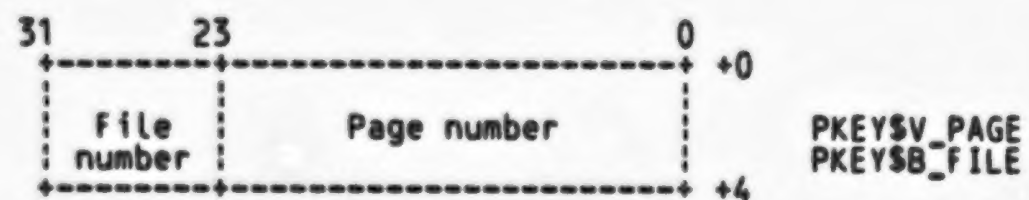
FIELD DKEYSZ_FIELDS =

SET

DKEYSV_LINE = [0, 0, 7, 0]
DKEYSV_PAGE = [0, 7, 17, 0]
DKEYSB_FILE = [0, 24, 8, 0]

TES;

In-core Page Number Structure



Dictionary page numbers that are passed around the hash table use the Page Key (PKEY) structure to provide both the page and file numbers.

LITERAL

PKEY\$S_BLOCK_LENGTH = 4;

MACRO

\$PKEY = BLOCK[PKEY\$S_BLOCK_LENGTH,BYTE] FIELD (PKEY\$Z_FIELDS)
%;

FIELD PKEY\$Z_FIELDS =

SET

PKEY\$V_PAGE
PKEY\$B_FILE

= [0, 0, 24, 0],
= [0, 24, 8, 0]

TES;

Line Index Format



Each block on a dictionary page is pointed to by a line index on that page. Disk pointers (DPTR) between the various blocks actually point to the block's line index. The line index is then used to locate the physical dictionary block.

The PAGE\$INDEX macro allows the programmer to access a line index while it resides on a dictionary page. It also uses the LINE\$Z_FIELDS to access the individual fields within a line index.

LITERAL

LINE\$S_BLOCK_LENGTH = 2;

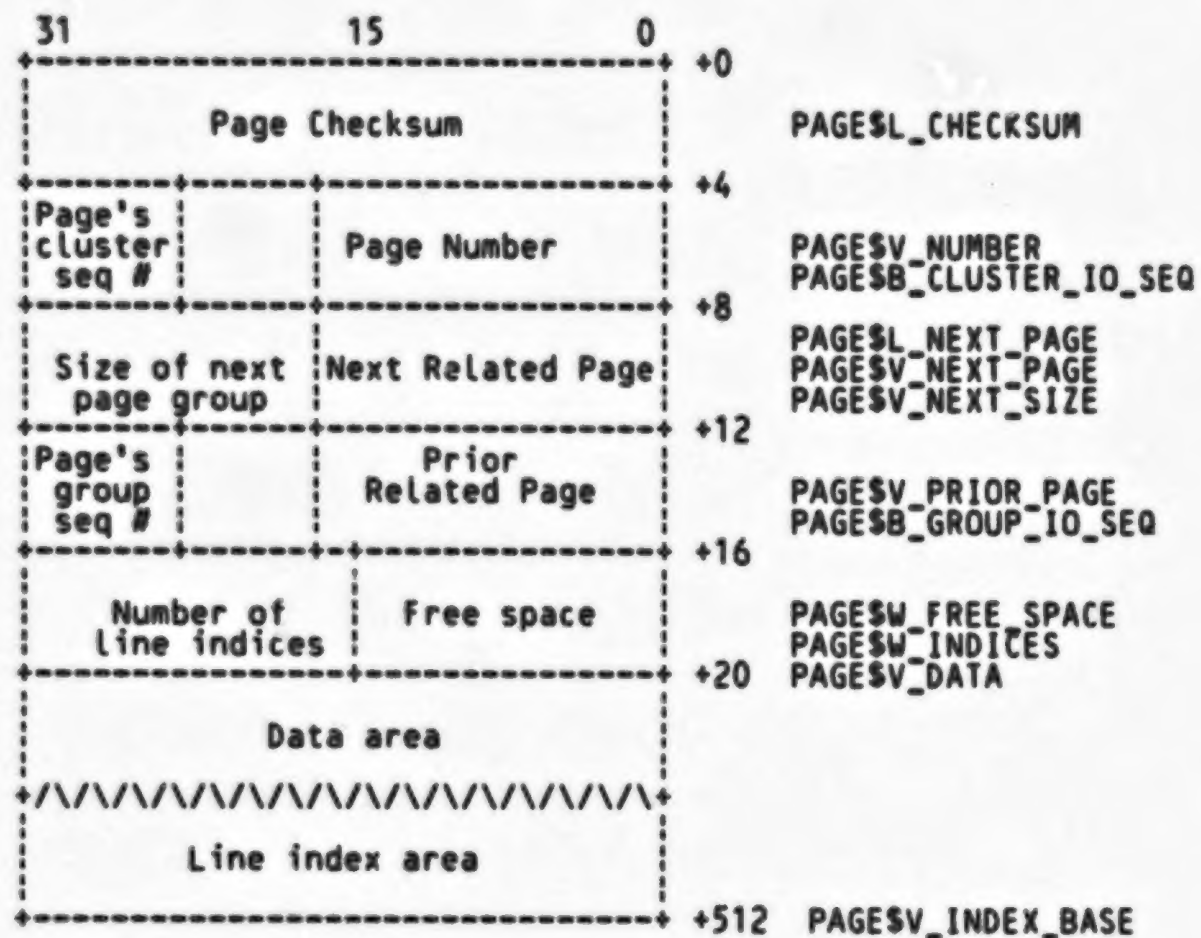
MACRO

\$LINE = BLOCK[LINE\$S_BLOCK_LENGTH, BYTE] FIELD (LINE\$Z_FIELDS)
%;

FIELD LINE\$Z_FIELDS =

SET
TES; LINE\$W_BLOCK_OFFSET = [0, 0, 16, 0]

Page Format



Each page in the dictionary file starts with a page header. The next and prior related page pointers are used for linking page groups. The following types of page groups exist:

- 1) Lockable page group.
Each named entity (and history list) owns exactly one lockable page group. These groups contain all unnamed offspring of the name entity (or history list).

A page group may only be accessed through its portal page. A group's portal page is the page on which the named entity (or history list head) resides. If a group's portal page is locked, then no pages in the group may be accessed.

The PAGESINDEX structure is used to access a line index entry on a page.

A page whose checksum is zero is a locked page, and indicates that

the sub-tree below it is incomplete and in a transient state.
Such pages may never be read.

Each page has two page sequence numbers. The cluster sequence number must be the same for all pages in the cluster. The group sequence number must be the same for all pages in an I/O group.

The cluster I/O sequence number is bumped whenever more than one group in the cluster is written. The group I/O sequence number is bumped whenever the group is written. These sequence numbers enable the detection of incomplete cluster unstage operations.

LITERAL

PAGESS_BLOCK_LENGTH = 512;

MACRO

\$PAGE = BLOCK[PAGESS_BLOCK_LENGTH, BYTE] FIELD (PAGE\$Z_FIELDS)
X;

FIELD PAGE\$Z_FIELDS =

SET

PAGE\$L_CHECKSUM	= [0, 0, 32, 0];
PAGE\$V_NUMBER	= [4, 0, 17, 0];
PAGE\$B_CLUSTER_IO_SEQ	= [7, 0, 8, 0];
PAGE\$L_NEXT_PAGE	= [8, 0, 32, 0];
PAGE\$V_NEXT_PAGE	= [8, 0, 17, 0];
PAGE\$V_NEXT_SIZE	= [8, 17, 15, 0];
PAGE\$V_PRIOR_PAGE	= [12, 0, 17, 0];
PAGE\$B_GROUP_IO_SEQ	= [15, 0, 8, 0];
PAGE\$W_FREE_SPACE	= [16, 0, 16, 0];
PAGE\$W_INDICES	= [18, 0, 16, 0];
PAGE\$V_DATA	= [20, 0, 0, 0];
PAGE\$V_INDEX_BASE	= [512, 0, 0, 0];

TES;

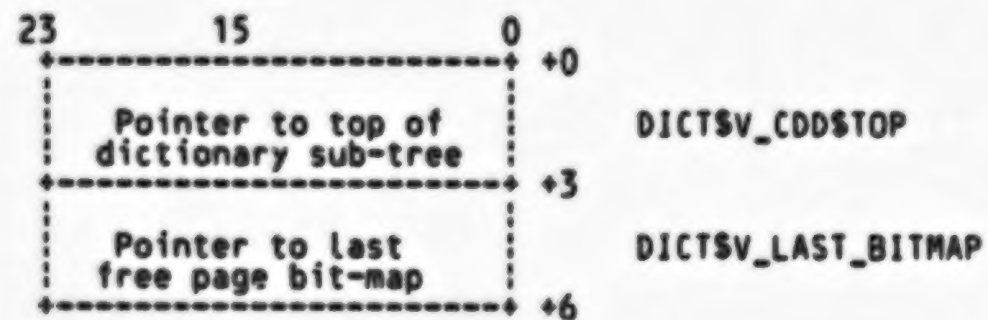
LITERAL

PAGE\$K_BASE = BLOCK[0, PAGE\$V_DATA;
PAGESS_BLOCK_LENGTH, BYTE];
PAGE\$K_FREE_SPACE = PAGESS_BLOCK_LENGTH - PAGE\$K_BASE;
PAGE\$S_INDEX_BASE = BLOCK[0, PAGE\$V_INDEX_BASE;
PAGESS_BLOCK_LENGTH, BYTE];

STRUCTURE

PAGE\$INDEX[I, 0, P, S, E] =
(PAGE\$INDEX+PAGE\$S_INDEX_BASE+0-((I)*LINESS_BLOCK_LENGTH)) <P,S,E>;

Dictionary Header Block



Each dictionary uses page 1 as its dictionary header page. The dictionary header block immediately follows the page header on page 1.

DICTSV_CDDSTOP points to the root of the tree/sub-tree contained in the dictionary file.

DICTSV_LAST_BITMAP points to the last free page bit-map entry.

LITERAL

DICTSS_BLOCK_LENGTH = 6 + PAGESK_BASE;

MACRO

SDICT = BLOCK[DICTSS_BLOCK_LENGTH, BYTE] FIELD (PAGESZ_FIELDS, DICTSZ_FIELDS)

%;

FIELD DICTSZ_FIELDS =

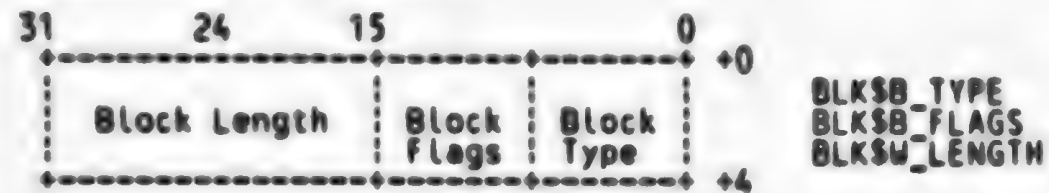
SET

DICTSV_CDDSTOP = [0+PAGESK_BASE, 0, 24, 0],

DICTSV_LAST_BITMAP = [3+PAGESK_BASE, 0, 24, 0]

TES;

Universal Block Header



This block header is used in every non-fixed block on the dictionary pages. Fixed blocks (one which always occur in the same place whenever they exist) do not have this block header.

LITERAL
BLK\$S_BLOCK_LENGTH = 4;

MACRO
\$BLK = BLOCK[BLK\$S_BLOCK_LENGTH,BYTE] FIELD (ACLSZ_FIELDS,
ACLC\$Z_FIELDS,
ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
ELST\$Z_FIELDS,
ENT\$Z_FIELDS,
FPCBSZ_FIELDS,
LIST\$Z_FIELDS,
LST\$Z_FIELDS,
NAME\$Z_FIELDS,
NAT\$Z_FIELDS,
NNAM\$Z_FIELDS,
NODE\$Z_FIELDS,
SEGSZ_FIELDS,
SLST\$Z_FIELDS,
SSASZ_FIELDS,
STR\$Z_FIELDS,
TEXT\$Z_FIELDS)

;

FIELD BLK\$Z_FIELDS =
SET
BLK\$B_TYPE = [0, 0, 8, 0],
BLK\$B_FLAGS = [1, 0, 8, 0],
BLK\$W_LENGTH = [2, 0, 16, 0]
TES;

LITERAL
BLK\$K_FLAGS = BLOCK[0, BLK\$B_TYPE; ,BYTE],
BLK\$K_BASE = BLK\$S_BLOCK_LENGTH;

!+

Block types

The ordering and clustering of these block types is significant.

All block types must be contiguous without any holes and bounded by the symbols BLKSK_TYPE_FIRST and BLKSK_TYPE_LAST.

All node and NAME block types must be contiguous and bounded by the symbols BLKSK_TYPE_NODE_FIRST and BLKSK_TYPE_NODE_LAST.

To add more node or NAME block types, insert them at the front of the table.

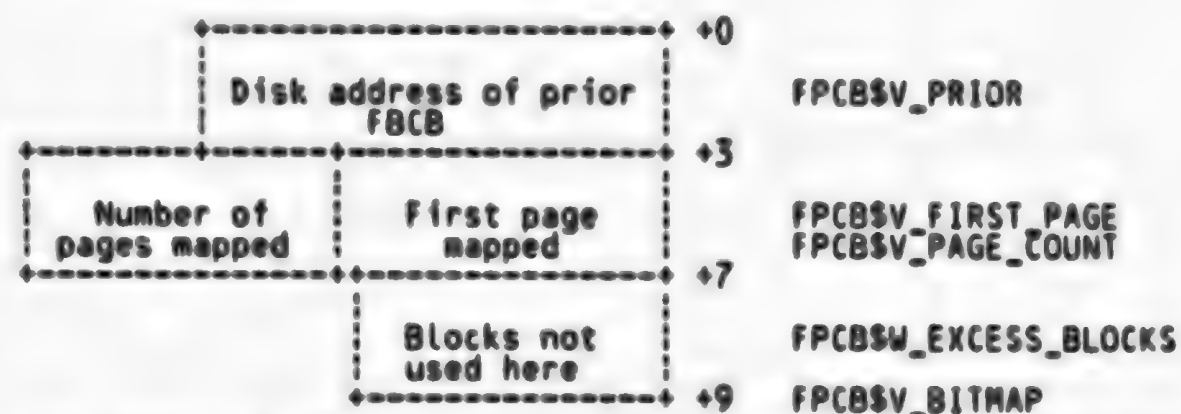
To add any other block type, append them to the end of the table.

CHANGING THE VALUES OF ANY OF THE ACTUAL BLOCK TYPES WILL INVALIDATE EXISTING DICTIONARIES.

LITERAL

BLKSK_TYPE_FIRST	= 101, ! Lowest block type
BLKSK_TYPE_NODE_FIRST	= 101, ! First node or NAM block
BLKSK_TYPE_DIR_NAM	= 101, ! Directory NAM block
BLKSK_TYPE_FIL_NAM	= 102, ! File NAM block
BLKSK_TYPE_TERM_NAM	= 103, ! Terminal NAM block
BLKSK_TYPE_DIR_NODE	= 104, ! Directory node block
BLKSK_TYPE_TERM_NODE	= 105, ! Terminal node block
BLKSK_TYPE_NODE_LAST	= 105, ! Last node or NAM block
BLKSK_TYPE_BITMAP	= 106, ! Free page bitmap
BLKSK_TYPE_ENTITY_ATT	= 107, ! Entity attribute block
BLKSK_TYPE_ENTITY_LIST	= 108, ! Entity list block
BLKSK_TYPE_ENTITY_LIST_ATT	= 109, ! Entity list attribute block
BLKSK_TYPE_NULL_ATT	= 110, ! Null attribute block
BLKSK_TYPE_NUM_ATT	= 111, ! Numeric attribute block
BLKSK_TYPE_SHORT_ATT	= 112, ! Short string attribute block
BLKSK_TYPE_STRING_ATT	= 113, ! String attribute block
BLKSK_TYPE_STRING_LIST	= 115, ! String list block
BLKSK_TYPE_STRING_LIST_ATT	= 116, ! String list attribute block
BLKSK_TYPE_STRING_SEG	= 117, ! String segment block
BLKSK_TYPE_TEXT	= 118, ! Text block
BLKSK_TYPE_ACL	= 119, ! Access control list entry
BLKSK_TYPE_ACLC	= 120, ! Access control list criterion
BLKSK_TYPE_LAST	= 120, ! Highest block type

Free Page Control Block (FPCB)
Free Page Bit Map (FPBM)



Free pages in the dictionary file are controlled by the Free Page Control Blocks (FPCB) and the Free Page Bit Map (FPBM). Each free page in the file has its corresponding bit set on.

The disk's cluster size may make it impossible to use all of the allocated blocks. If so, the un-used blocks are tallied in FPCBSW_EXCESS_BLOCKS, and will be used when the file is next extended.

The FPCBs are chained as the dictionary file grows in size.

LITERAL

FPCBS_BLOCK_LENGTH = BLKSK_BASE + 9;

STRUCTURE

FPCB[0, P, S, E; BLOCKS] = [FPCBS_BLOCK_LENGTH+(BLOCKS+7)/8]
(FPCB+0)<P,S,E>;

FIELD FPCBSZ_FIELDS =

SET

FPCBSV_PRIOR	= [0+BLKSK_BASE, 0, 24, 0],
FPCBSV_FIRST_PAGE	= [3+BLKSK_BASE, 0, 17, 0],
FPCBSV_PAGE_COUNT	= [3+BLKSK_BASE, 17, 15, 0],
FPCBSW_EXCESS_BLOCKS	= [7+BLKSK_BASE, 0, 16, 0],
FPCBSV_BITMAP	= [9+BLKSK_BASE, 0, 0, 0]

TES;

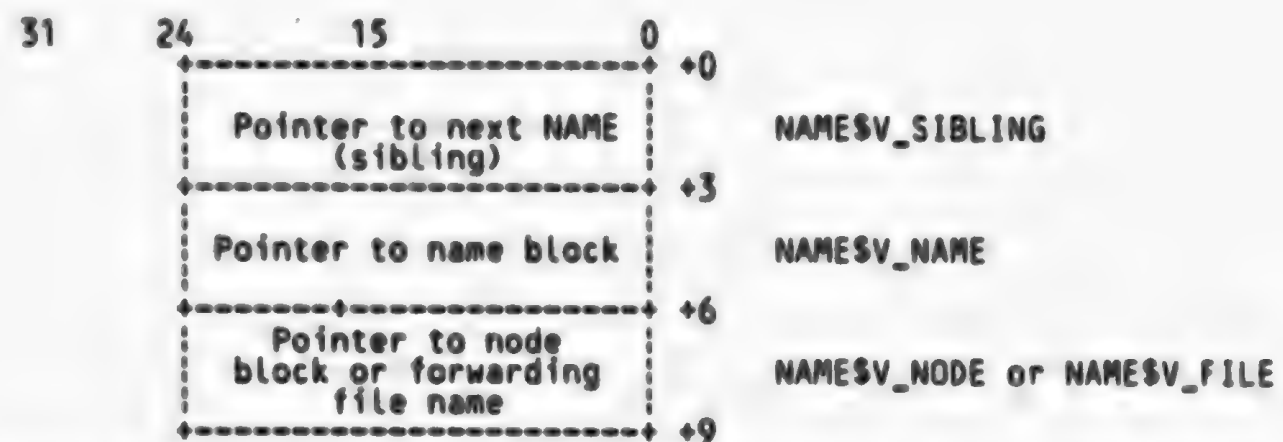
LITERAL

FPCBSK_BASE = FPCBS_BLOCK_LENGTH,
FPCBS_BITMAP = FPCB[0, FPCBSV_BITMAP];

STRUCTURE

FPBM[1] = (FPBM+FPCBS_BITMAP)<1-1, 1, 0>;

Name Block



There are three types of NAM Blocks, Directory, File, and Terminal.

Directory NAM Block is made up of a Block Header + NAM Block.

File NAM Block is made up of a Block Header + NAM Block.

Terminal NAM Block is made up of a Block Header + NAM Block + Terminal NAM Block.

LITERAL

NAMES_BLOCK_LENGTH = 9 + BLKSK_BASE;

MACRO

\$NAME = BLOCK[NAMES_BLOCK_LENGTH, BYTE] FIELD (BLKSZ_FIELDS, NAMESZ_FIELDS)

%;

FIELD NAMESZ_FIELDS =

SET

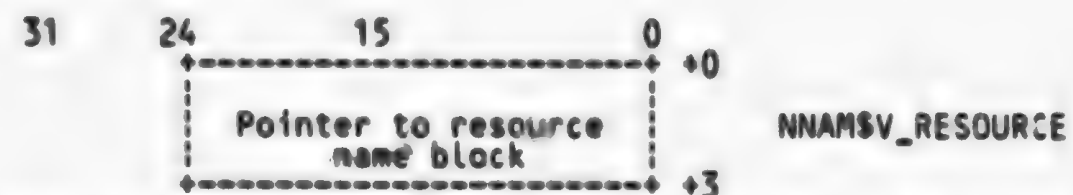
NAMESV_SIBLING	= [0+BLKSK_BASE, 0, 24, 0],
NAMESV_NAME	= [3+BLKSK_BASE, 0, 24, 0],
NAMESV_NODE	= [6+BLKSK_BASE, 0, 24, 0],
NAMESV_FILE	= [6+BLKSK_BASE, 0, 24, 0]

TES,

LITERAL

NAMESK_BASE = NAMES_BLOCK_LENGTH;

Node Nam Block



Node NAM Block is made up of a Block Header + NAM Block + Node NAM Block.

LITERAL

NNAMSS_BLOCK_LENGTH = 3 + NAME\$K_BASE;

MACRO

\$NNAM = BLOCK[NNAMSS_BLOCK_LENGTH,BYTE] FIELD (BLK\$Z_FIELDS,
NAME\$Z_FIELDS,
NNAM\$Z_FIELDS)

%;

FIELD NNAM\$Z_FIELDS =

SET

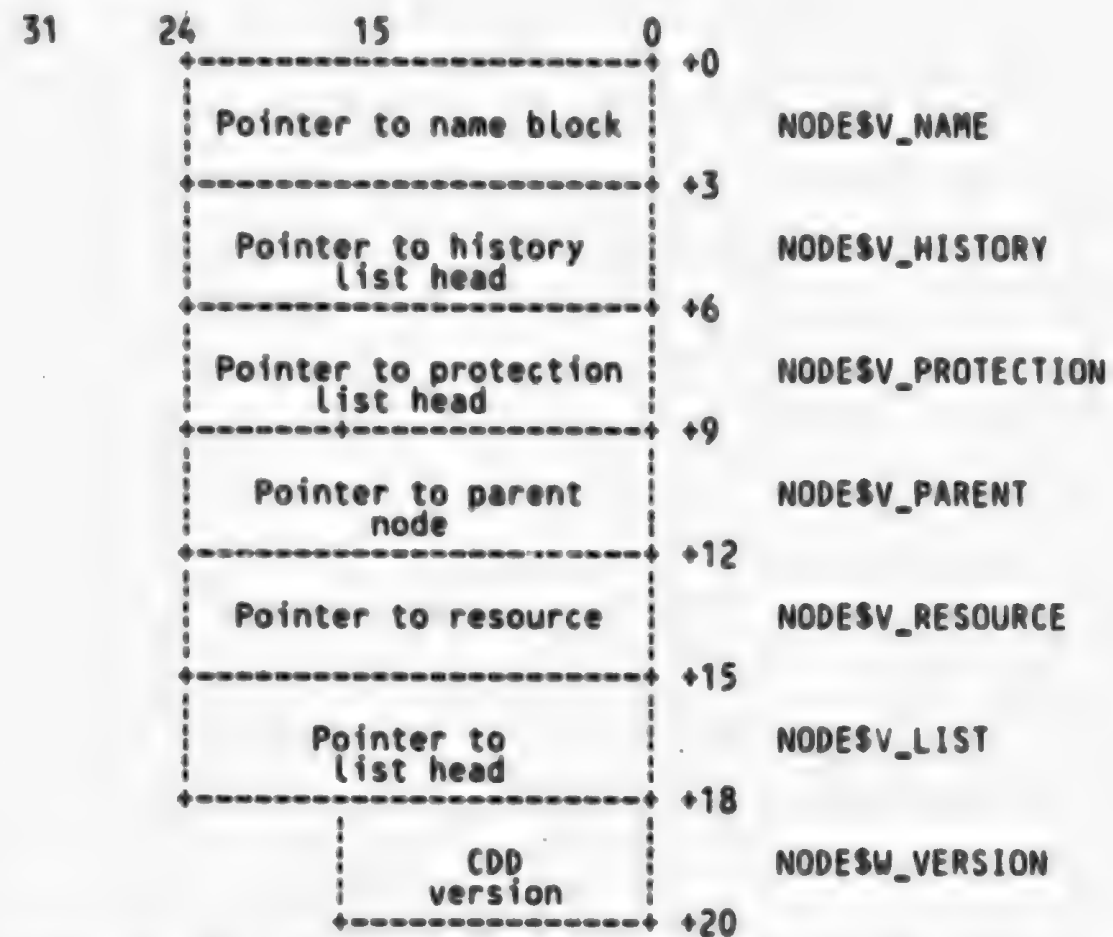
NNAMSV_RESOURCE = [0+NAME\$K_BASE, 0, 24, 0]

TES;

LITERAL

NNAM\$K_BASE = NNAMSS_BLOCK_LENGTH;

Node Block



There are two types of Node Blocks Directory and Terminal.
 Directory Node Block is made up of a Block Header + Node Block
 Terminal Node Block is made up of a Block Header + Node Block

LITERAL

NODE\$S_BLOCK_LENGTH = 20+ BLK\$K_BASE;

MACRO

\$NODE = BLOCK[NODE\$S_BLOCK_LENGTH,BYTE] FIELD (BLK\$Z_FIELDS,
 NODE\$Z_FIELDS)

1;

FIELD NODE\$Z_FIELDS =

SET

NODE\$V_ORDERED = [1, 0, 1, 0]
 NODE\$V_NAME = [0+BLK\$K_BASE, 0, 24, 0].
 NODE\$V_HISTORY = [3+BLK\$K_BASE, 0, 24, 0].

NODESV_PROTECTION	=	[6+BLK\$K_BASE, 0, 24, 0],
NODESV_PARENT	=	[9+BLK\$K_BASE, 0, 24, 0],
NODESV_RESOURCE	=	[12+BLK\$K_BASE, 0, 24, 0],
NODESV_LIST	=	[15+BLK\$K_BASE, 0, 24, 0],
NODESM_VERSION	=	[18+BLK\$K_BASE, 0, 16, 0]

TES;

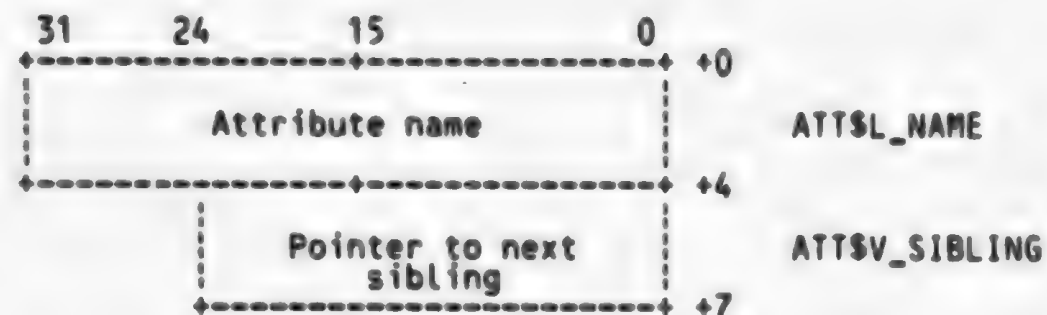
LITERAL
NODE\$K_BASE = NODE\$S_BLOCK_LENGTH;

|+
|
|-

The following flag occurs in the BLK\$B_FLAGS field of
a directory node.

LITERAL
NODE\$M_ORDERED = 1^1 - 1^0; ! List elements are sorted by name

Common Attribute Block



Each attribute has the common attribute block immediately following the universal block header.

There are six types of Attribute Blocks Entity Attribute Block, List Attribute Block, Numeric Attribute Block, Null Attribute Block, Short String Attribute Block, and String Attribute Block.

Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.

List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.

Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.

Null Attribute Block is made up of a Block Header + Attribute Block.

Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

LITERAL

ATT\$S_BLOCK_LENGTH = 7 + BLK\$K_BASE;

MACRO

\$ATT = BLOCK[ATT\$S_BLOCK_LENGTH, BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS)

;

FIELD ATT\$Z_FIELDS =

SET

ATTSL_NAME = [0+BLK\$K_BASE, 0, 32, 0],
ATT\$V_SIBLING = [4+BLK\$K_BASE, 0, 24, 0]

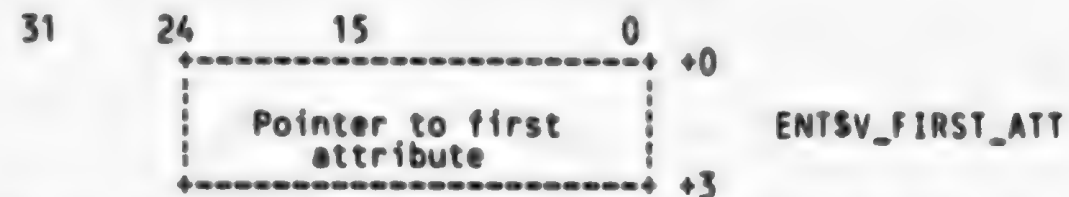
TES;

LITERAL

ATT&K_BASE

= ATT&S_BLOCK_LENGTH;

Entity Attribute Block



Entity Attribute Block is made up of a Block Header + Attribute Block
+ Entity Attribute Block.

LITERAL

ENT\$S_BLOCK_LENGTH = 3 + ATT\$K_BASE;

MACRO

SENT = BLOCK[ENT\$S_BLOCK_LENGTH, BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
ENT\$Z_FIELDS)

%;

FIELD ENT\$Z_FIELDS =

SET

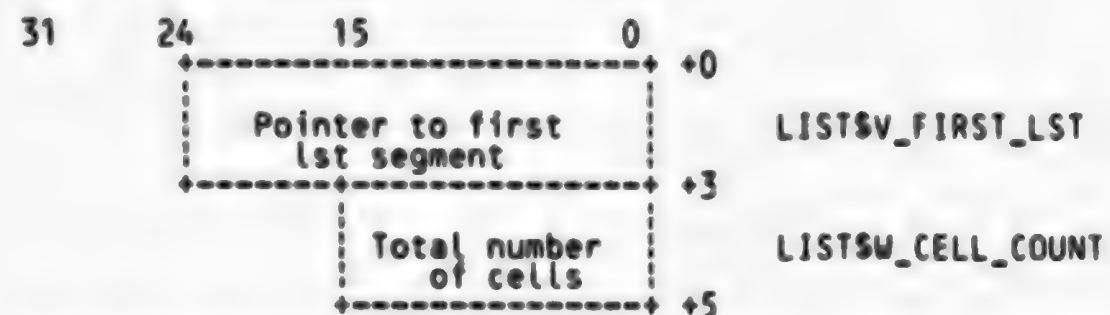
ENTSV_FIRST_ATT = [0+ATT\$K_BASE, 0, 24, 0]

TES;

LITERAL

ENT\$K_BASE = ENT\$S_BLOCK_LENGTH;

List Attribute Block



List Attribute Block is made up of a Block Header + Attribute Block
+ List Attribute Block.

LITERAL

LIST\$\$BLOCK_LENGTH = 5 + ATT\$K_BASE;

MACRO

\$LIST = BLOCK[LIST\$\$BLOCK_LENGTH, BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
LIST\$Z_FIELDS)

%;

FIELD LIST\$Z_FIELDS =

SET

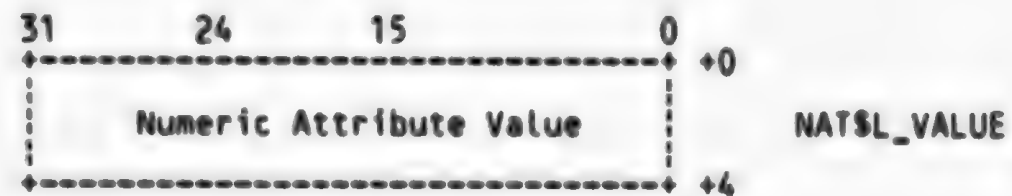
LIST\$V_FIRST_LST = [0+ATT\$K_BASE, 0, 24, 0],
LIST\$W_CELL_COUNT = [3+ATT\$K_BASE, 0, 16, 0]

TES;

LITERAL

LIST\$K_BASE = LIST\$\$BLOCK_LENGTH;

Numeric Attribute Block



Numeric Attribute Block is made up of a Block Header + Attribute Block
+ Numeric Attribute Block.

Numeric attribute block contains the
value of the numeric attribute.

```
LITERAL
  NAT$S_BLOCK_LENGTH = 4 + ATT$K_BASE;
```

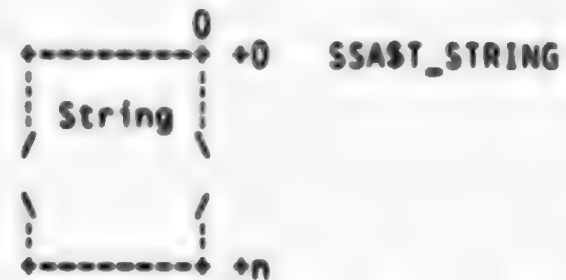
```
MACRO
  $NAT = BLOCK[NAT$S_BLOCK_LENGTH,BYTE] FIELD (ATT$Z_FIELDS,
                                                BLK$Z_FIELDS,
                                                NAT$Z_FIELDS)
  X;
```

```
FIELD NAT$Z_FIELDS =
  SET
    NAT$L_VALUE          = [0+ATT$K_BASE,0,32,0]
  TES;
```

```
LITERAL
  NAT$K_BASE            = NAT$S_BLOCK_LENGTH;
```

Short String Attribute Block

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Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

Strings whose total length is between 0 and 255 bytes are usually stored using the short string attribute block. If a string is too long to be stored in an SSA, then it is stored using a normal string attribute block (STR).

LITERAL

SSASS_BLOCK_LENGTH = 0 + ATT\$K_BASE;

MACRO

\$SSA = BLOCK[SSASS_BLOCK_LENGTH, BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
SSASZ_FIELDS)

%;

FIELD SSASZ_FIELDS =

SET

SSAST_STRING

YES;

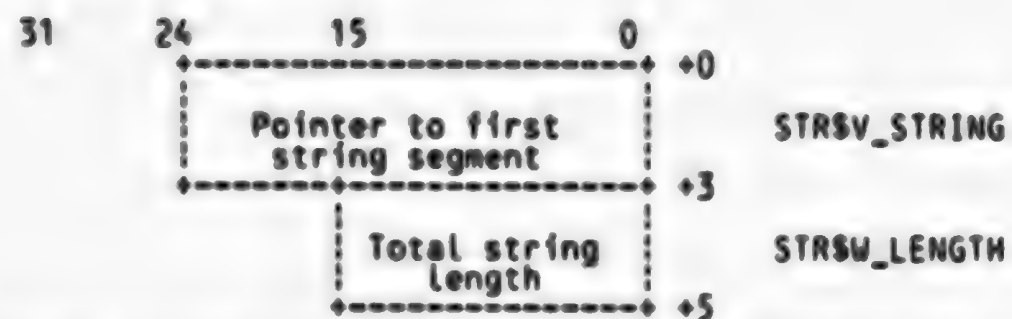
= [0+ATT\$K_BASE, 0, 0, 0]

LITERAL

SSASK_BASE

= SSASS_BLOCK_LENGTH;

String Attribute Block



String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

The string is broken into one or more segments. The string segments are stored in SEG blocks.

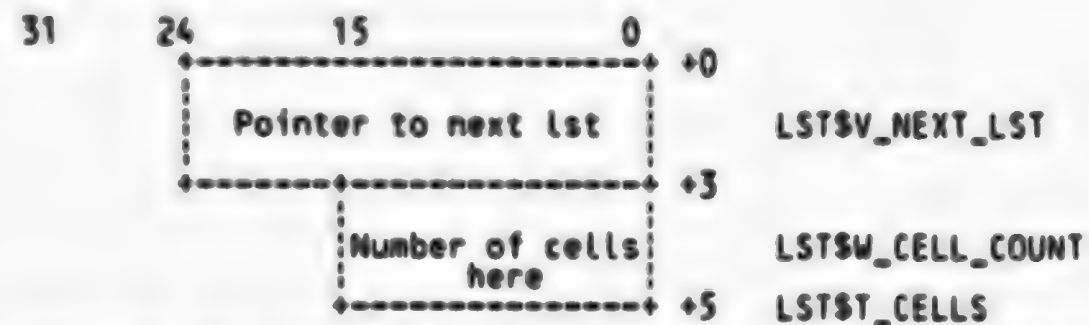
LITERAL
STRSS_BLOCK_LENGTH = 5 + ATT\$K_BASE;

MACRO
\$STR = BLOCK[STRSS_BLOCK_LENGTH, BYTE] FIELD (ATT\$Z_FIELDS,
BLK\$Z_FIELDS,
STR\$Z_FIELDS)
%;

FIELD STR\$Z_FIELDS =
SET
STRSV_STRING = [0+ATT\$K_BASE, 0, 24, 0],
STRSW_LENGTH = [3+ATT\$K_BASE, 0, 16, 0]
YES;

LITERAL
STR\$K_BASE = STRSS_BLOCK_LENGTH;

LST Segment Block



There are two types of LST Segment Blocks Entity List Block and String List Attribute Block.

Entity List Block is made up of a Block Header + LST Block + Entity List Block.

String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

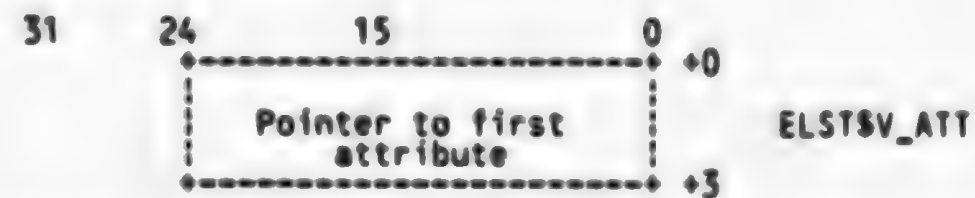
LITERAL
LSTSS_BLOCK_LENGTH = 5 + BLKSK_BASE;

MACRO
\$LST = BLOCK[LSTSS_BLOCK_LENGTH, BYTE] FIELD (BLKSZ_FIELDS,
LSTSZ_FIELDS)
%;

FIELD LSTSZ_FIELDS =
SET
LSTSV_NEXT_LST = [0 + BLKSK_BASE, 0, 24, 0],
LSTSW_CELL_COUNT = [3 + BLKSK_BASE, 0, 16, 0],
LSTST_CELLS = [5 + BLKSK_BASE, 0, 0, 0]
TES;

LITERAL
LSTSK_BASE = LSTSS_BLOCK_LENGTH;

Entity List Block



Entity List Block is made up of a Block Header + LST Block + Entity List Block.

```
LITERAL
  ELSTSS_BLOCK_LENGTH = 3;
```

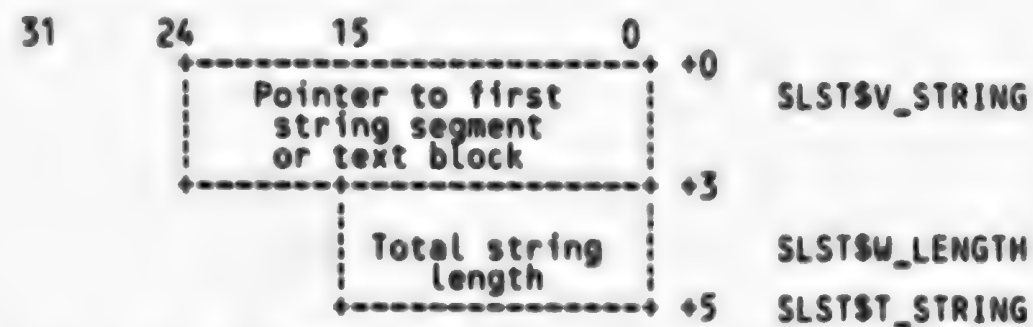
```
MACRO
  $ELST = BLOCK[ELSTSS_BLOCK_LENGTH, BYTE] FIELD (ELST$Z_FIELDS)
  %;
```

```
FIELD ELST$Z_FIELDS =
  SET
    ELSTSV_ATT      = [0, 0, 24, 0]
  TES;
```

```
LITERAL
  ELST$K_BASE      = ELSTSS_BLOCK_LENGTH;
```

```
STRUCTURE
  ELST$ELM[I] =
    (ELST$ELM + LST$K_BASE + (I * ELSTSS_BLOCK_LENGTH));
```

String List Attribute Block



String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

```
LITERAL
  SLST$S_BLOCK_LENGTH = 5;
```

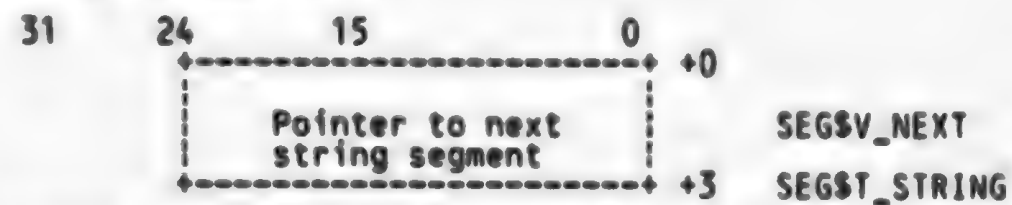
```
MACRO
  $SLST = BLOCK[SLST$S_BLOCK_LENGTH, BYTE] FIELD (SLST$Z_FIELDS)
  %;
```

```
FIELD SLST$Z_FIELDS =
  SET
    SLST$V_STRING      = [0, 0, 24, 0];
    SLST$W_LENGTH      = [3, 0, 16, 0];
    SLST$T_STRING      = [5, 0, 0, 0];
  TES;
```

```
LITERAL
  SLST$K_BASE          = SLST$S_BLOCK_LENGTH;
```

```
STRUCTURE
  SLST$ELM[I] =
    (SLST$ELM + LST$K_BASE + (I * SLST$S_BLOCK_LENGTH));
```

String Segment Block



String Segment Block is made up of a Block Header + String Segment Block.

Each string attribute points to zero or more string segment blocks. Each block contains a portion of the whole string. The final string is found by concatenating all the string segments together.

```
LITERAL
SEG$$BLOCK_LENGTH = 3 + BLK$$BASE;
```

```
MACRO
    $SEG = BLOCK[SEG$$_BLOCK_LENGTH,BYTE] FIELD (BLK$$_FIELDS,
    $:                                     SEG$$_FIELDS)
```

```

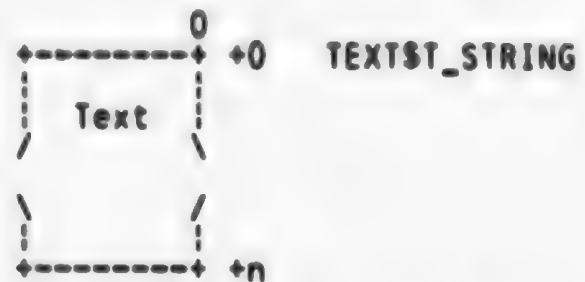
FIELD SEG$Z_FIELDS =
  SET
    SEG$V_NEXT           = [0+BLK$K_BASE, 0, 24, 0],
    SEG$T_STRING         = [3+BLK$K_BASE, 0, 0, 0]
  TES;

```

```
LITERAL
  SEGSS_BASE      = SEGSS_BLOCK_LENGTH:
```

Text Block

31



Text Block is made up of a Block Header + Text Block.

LITERAL

TEXT\$S_BLOCK_LENGTH = 0 + BLK\$K_BASE;

MACRO

```
$TEXT = BLOCK[TEXT$S_BLOCK_LENGTH,BYTE] FIELD (BLK$Z_FIELDS,
                                                TEXT$Z_FIELDS)
```

%;

FIELD TEXT\$Z_FIELDS =

SET

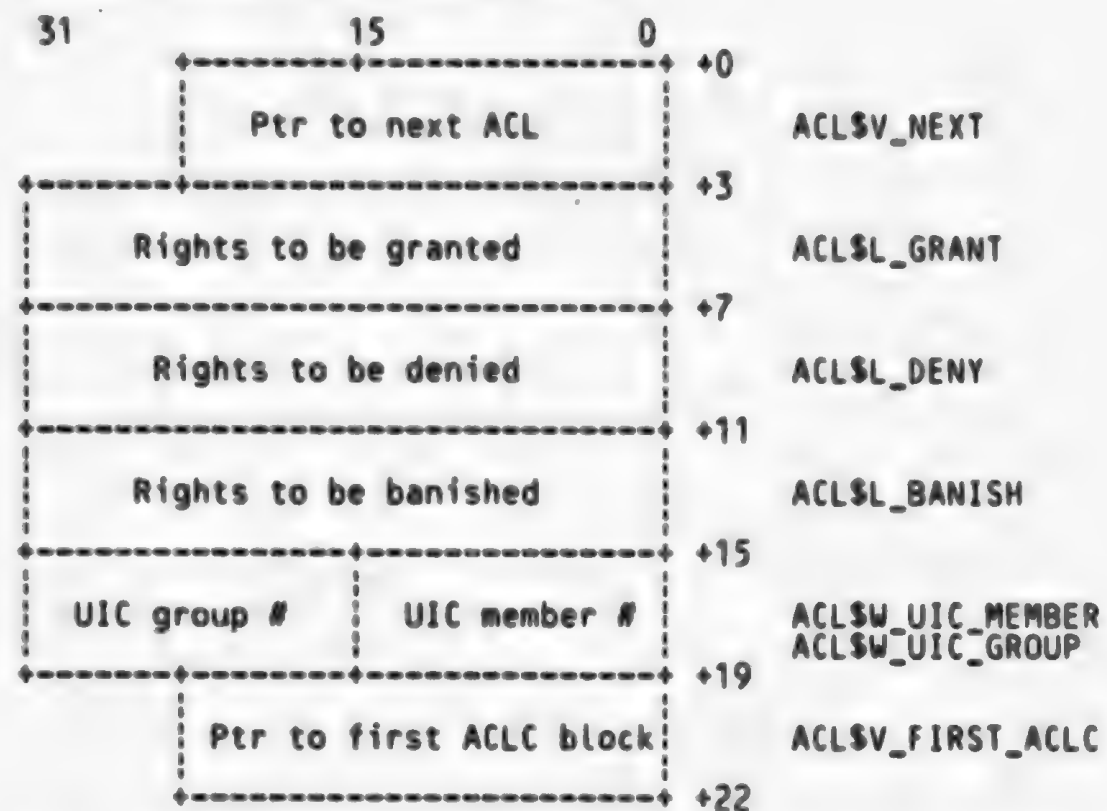
```
TEXT$T_STRING . = [0+BLK$K_BASE, 0, 0, 0]
```

TES;

LITERAL

TEXT\$K_BASE = TEXT\$S_BLOCK_LENGTH;

Access Control List Entry (ACL)



Each node has an Access Control List made up of zero or more Access Control List Entries (ACL).

The ACL block is appended to a BLK to make an access control list entry.

LITERAL

ACLS_BLOCK_LENGTH = 22+BLK\$K_BASE;

MACRO

\$ACL = BLOCK[ACLS_BLOCK_LENGTH, BYTE] FIELD (BLK\$Z_FIELDS,
ACLS\$Z_FIELDS)

%;

FIELD SET ACL\$Z_FIELDS =

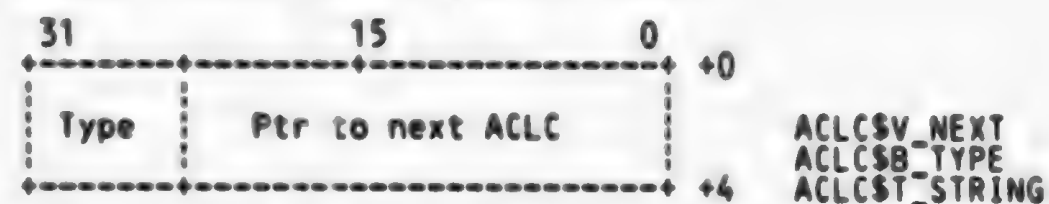
ACLSV_NEXT	= [0+BLK\$K_BASE, 0, 24, 0],
ACLSL_GRANT	= [3+BLK\$K_BASE, 0, 32, 0],
ACLSL_DENY	= [7+BLK\$K_BASE, 0, 32, 0],
ACLSL_BANISH	= [11+BLK\$K_BASE, 0, 32, 0],
ACLSW_UIC_MEMBER	= [15+BLK\$K_BASE, 0, 16, 0],
ACLSW_UIC_GROUP	= [17+BLK\$K_BASE, 0, 16, 0],
ACLSV_FIRST_ACLC	= [19+BLK\$K_BASE, 0, 24, 0]

TES;

LITERAL
ACL\$K_BASE

= ACL\$S_BLOCK_LENGTH;

Access Control List Criterion (ACLC)



Each ACL may have one or more Access Control List Criterion (ACLC) chained from it.

Each ACLC specifies one user identification criterion for the ACL entry. The user must match all the criteria for the ACL entry to apply to him.

An ACLC is appended to a BLK to form the criterion block.

LITERAL

ACLC\$S_BLOCK_LENGTH = 4+BLK\$K_BASE;

MACRO

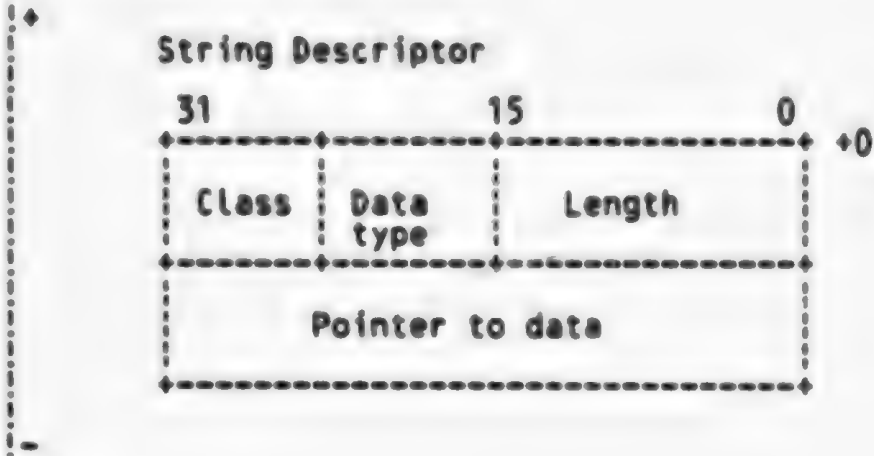
\$ACLC = BLOCK[ACLC\$S_BLOCK_LENGTH, BYTE] FIELD (BLK\$Z_FIELDS,
ACLC\$Z_FIELDS)
%;

FIELD ACLC\$Z_FIELDS =

SET
ACLC\$V_NEXT = [0+BLK\$K_BASE, 0, 24, 0],
ACLC\$B_TYPE = [3+BLK\$K_BASE, 0, 8, 0],
ACLC\$T_STRING = [4+BLK\$K_BASE, 0, 0, 0]
TES;

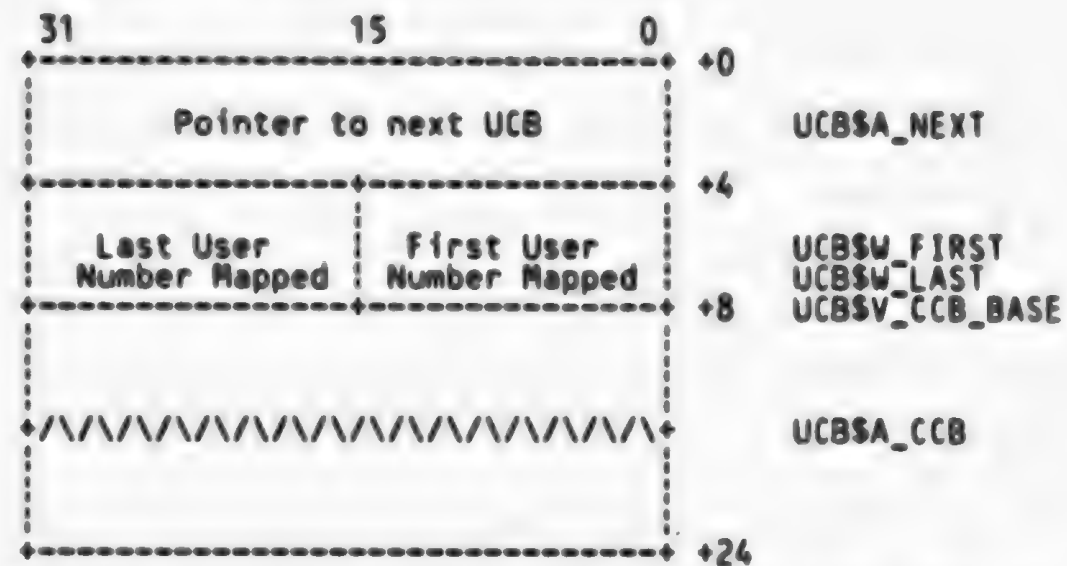
LITERAL

ACLC\$K_BASE = ACLC\$S_BLOCK_LENGTH;



```
MACRO
SDSC = BLOCK[8,BYTE]
%;
```

User Control Block (UCB)



A user's context pointer is an origin 1 index into the CCBs in the UCB list. Each UCB can point to UCB\$K_CCB number of CCBs. If the user passes a context number that doesn't map to an active CCB, we tell him it's an invalid context pointer.

Except in some extreme case, any image will probably never have more than 4 users active at any one time. But we can handle more!

LITERAL

```
UCB$K_CCB_BASE      = 8,
UCB$K_CCB           = 4,      ! Number of CCBs/UCB
UCB$S_BLOCK_LENGTH  = UCB$K_CCB_BASE + UCB$K_CCB * 4;
```

MACRO

```
$UCB = BLOCK[UCB$S_BLOCK_LENGTH, BYTE] FIELD (UCB$Z_FIELDS)
$;
```

FIELD UCB\$Z_FIELDS =

```
SET
  UCB$A_NEXT      = [0, 0, 32, 0],
  UCB$W_FIRST     = [4, 0, 16, 0],
  UCB$W_LAST      = [6, 0, 16, 0],
  UCB$V_CCB_BASE  = [8, 0, 0, 0],
  UCB$A_CCB       = [0, 0, 32, 0]
TES;
```

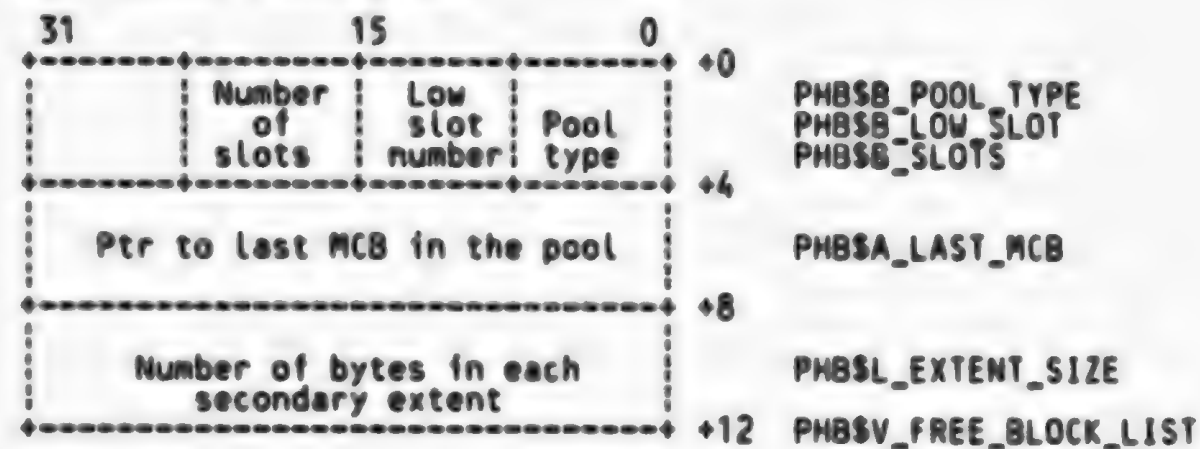
LITERAL

```
UCB$C_FIRST = BLOCK[0, UCB$W_FIRST; UCB$S_BLOCK_LENGTH, BYTE];
```

STRUCTURE

```
UCB$CCB[I, 0, P, S, E] = (UCB$CCB+UCB$K_CCB_BASE+
  (I-(UCB$CCB+UCB$C_FIRST)<0,16,0>)*4+0)<P,S,E>;
```

Pool Header Block (PHB)



Pools are used to provide space for various in-core block types. Blocks that are related are usually allocated from the same pool. This is done because pools provide good locality of reference and this allocation scheme reduces page faults.

Each pool consists of one or more extents. Each extent consists of a Memory Control Block and the rest of the space allocated to the extent. Each pool has a Pool Header Block associated with it. This block identifies the MCB list, as well as contains the pool's free block list.

The free block list is a vector of linked lists. Each slot in the vector corresponds to a block type. When a block is freed, it is linked into its associated free list. When a block is requested, its free list is checked to see if it is non-empty. If so, then a block is allocated from it. Otherwise, a block is allocated from one of the pool's extents.

```

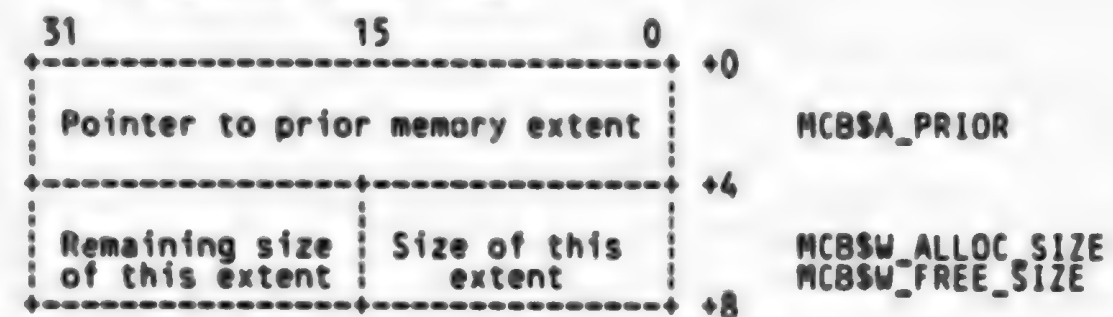
LITERAL
  PHBS_BLOCK_LENGTH      = 12;

MACRO
  $PHB = BLOCK[PHBS_BLOCK_LENGTH, BYTE] FIELD (PHBSZ_FIELDS)
  %;

FIELD PHBSZ_FIELDS =
  SET
    PHBSB_POOL_TYPE      = [0, 0, 8, 0],
    PHBSB_LOW_SLOT       = [1, 0, 8, 0],
    PHBSB_SLOTS          = [2, 0, 8, 0],
    PHBSA_LAST_MCB       = [4, 0, 32, 0],
    PHBSL_EXTENT_SIZE    = [8, 0, 32, 0],
    PHBSV_FREE_BLOCK_LIST = [12, 0, 0, 0]
  TES;

```

Memory Control Block (MCB)



Each pool has one or more Memory Control Blocks associated with it.

This list serves two purposes:

- 1) It points to each memory extent we asked LIB\$GET_VM for.
- 2) It keeps track of memory that has been allocated for the pool, but never used.

LITERAL

MCB\$BLOCK_LENGTH = 8;

MACRO

\$MCB = BLOCK[MCB\$BLOCK_LENGTH,BYTE] FIELD (MCB\$Z_FIELDS)
%;

FIELD

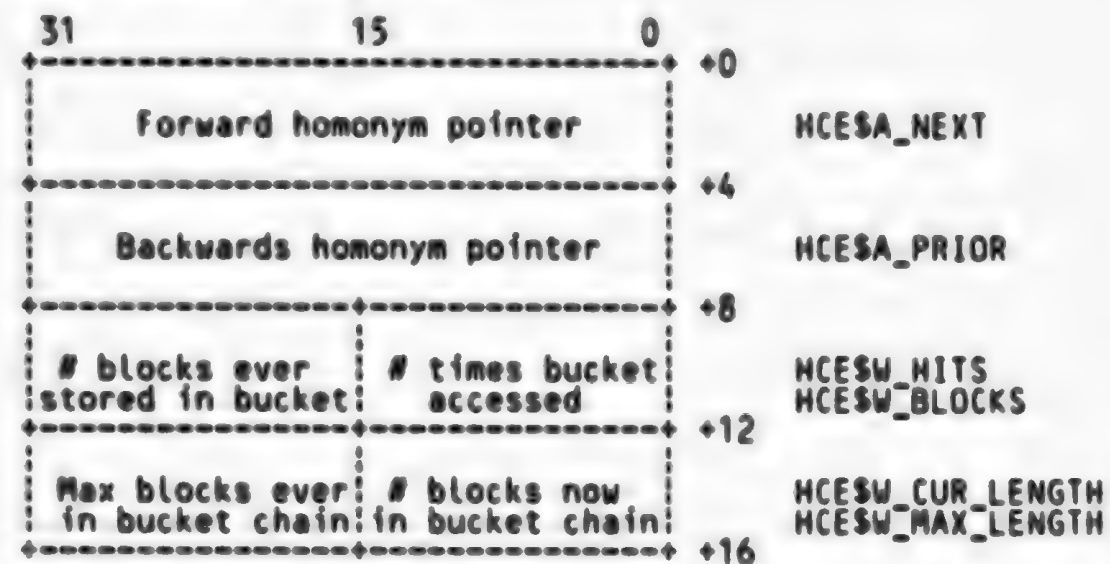
MCB\$Z_FIELDS =

SET

MCBSA_PRIOR = [0, 0, 32, 0],
MCBSW_ALLOC_SIZE = [4, 0, 16, 0],
MCBSW_FREE_SIZE = [6, 0, 16, 0]

TES;

Hash Control Entry (HCE)



HCEs are found in each user's CCB. Each HCE forms one hash bucket. We keep some statistics on the buckets in an attempt to tune the hash function.

LITERAL

HCESS_BLOCK_LENGTH = 16;

MACRO

\$HCE = BLOCK[HCESS_BLOCK_LENGTH, BYTE] FIELD (HCE\$Z_FIELDS)
%;

FIELD HCE\$Z_FIELDS =

SET

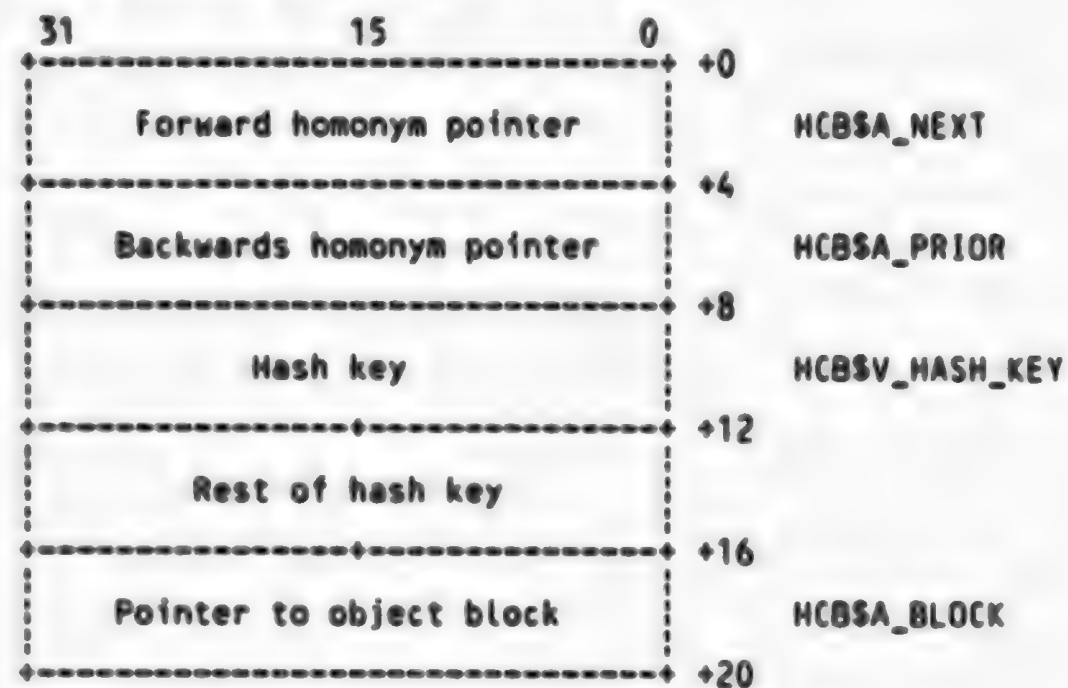
HCE\$A_NEXT	= [0, 0, 32, 0].
HCE\$A_PRIOR	= [4, 0, 32, 0].
HCE\$W_HITS	= [8, 0, 16, 0].
HCE\$W_BLOCKS	= [10, 0, 16, 0].
HCE\$W_CUR_LENGTH	= [12, 0, 16, 0].
HCE\$W_MAX_LENGTH	= [14, 0, 16, 0].

TES;

LITERAL

HCE\$K_HOMONYM_LIST = BLOCK[0, HCE\$A_NEXT; , BYTE];

Hash Control Block (HCB)



LITERAL

```

HCB$K_PCB_NUMBER      = 1.      ! Hash type is page number
HCB$K_LCCB_CODE       = 2.      ! Hash type is location code
HCB$K_LCCB_ADDRESS    = 3.      ! Hash type is entity disk address
HCB$K_KEY_LENGTH      = 8.      ! Size of a hash key
HCB$S_BLOCK_LENGTH    = 20;

```

MACRO

```

$HCB = BLOCK[HCB$S_BLOCK_LENGTH,BYTE] FIELD (HCB$Z_FIELDS)
%:

```

FIELD HCB\$Z_FIELDS =

```

SET
  HCB$A_NEXT      = [0, 0, 32, 0].
  HCB$A_PRIOR     = [4, 0, 32, 0].
  HCB$V_HASH_KEY  = [8, 0, 0, 0].
  HCB$A_BLOCK     = [16, 0, 32, 0].
TES;

```

LITERAL

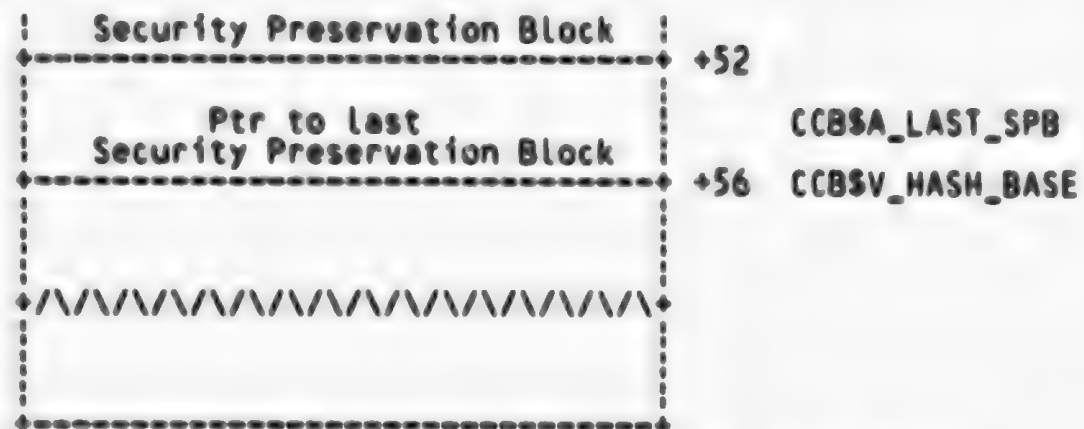
```

HCB$K_HOMONYM_LIST = BLOCK[0, HCB$A_NEXT; , BYTE];

```

Context Control Block (CCB)

31	15	0	+0	
Flags				CCBSW_FLAGS
Pointer to first LCB in temporary lock list			+4	CCBSA_FIRST_LOCK
Pointer to last LCB in temporary lock list			+8	CCBSA_LAST_LOCK
User ID			+12	CCBSW_LOCK_ID
User ID for lock manager			+16	CCBSW_USER_ID
Next location code number to be assigned			+20	CCBSL_NEXT_LCC
Ptr to NCB for file's CDDSTOP node			+24	CCBSA_CDDSTOP_NCB
NCB of CDD\$LOGIN node			+28	CCBSA_LOGIN_NCB
Ptr to current FCB for user			+32	CCBSA_CURRENT_FCB
Ptr to highest cluster in cache			+36	CCBSA_FIRST_CIB
Ptr to last highest cluster in cache			+40	CCBSA_LAST_CIB
Status of current transaction			+44	CCBSL_STATUS
Ptr to CCB's pool header			+48	CCBSA_POOL
Ptr to first				CCBSA_FIRST_SPB



The hash table consists of a number of hash entries.
See the HCE description for the format of these entries.

LITERAL

```
CCBSK_HASH_TABLE_SIZE = 151,
CCBS_BLOCK_LENGTH = 56 + (CCBSK_HASH_TABLE_SIZE * HCESS_BLOCK_LENGTH);
```

MACRO

```
$CCB = BLOCK[CCBS_BLOCK_LENGTH, BYTE] FIELD (CCBSZ_FIELDS)
%;
```

FIELD CCBSZ_FIELDS =

```
SET
  CCBSW_FLAGS = [0, 0, 16, 0],
  CCBSV_CORRUPT = [0, 0, 1, 0], ! Stream is corrupt
  CCBSA_FIRST_LOCK = [4, 0, 32, 0],
  CCBSA_LAST_LOCK = [8, 0, 32, 0],
  CCBSW_LOCK_ID = [12, 0, 16, 0],
  CCBSW_USER_ID = [14, 0, 16, 0],
  CCBSL_NEXT_LCC = [16, 0, 32, 0],
  CCBSA_CDDSTOP_NCB = [20, 0, 32, 0],
  CCBSA_LOGIN_NCB = [24, 0, 32, 0],
  CCBSA_CURRENT_FCB = [28, 0, 32, 0],
  CCBSA_FIRST_CIB = [32, 0, 32, 0],
  CCBSA_LAST_CIB = [36, 0, 32, 0],
  CCBSL_STATUS = [40, 0, 32, 0],
  CCBSA_POOL = [44, 0, 32, 0],
  CCBSA_FIRST_SPB = [48, 0, 32, 0],
  CCBSA_LAST_SPB = [52, 0, 32, 0],
  CCBSV_HASH_BASE = [56, 0, 0, 0]
TES;
```

LITERAL

```
CCBSK_LOCK_LIST = BLOCK[0, CCBSA_FIRST_LOCK, , BYTE],
CCBSK_CLUSTER_LIST = BLOCK[0, CCBSA_FIRST_CIB, , BYTE],
CCBSV_HASH_BASE = BLOCK[0, CCBSV_HASH_BASE, , BYTE],
CCBSK_SPB_LIST = BLOCK[0, CCBSA_FIRST_SPB, , BYTE];
```

STRUCTURE

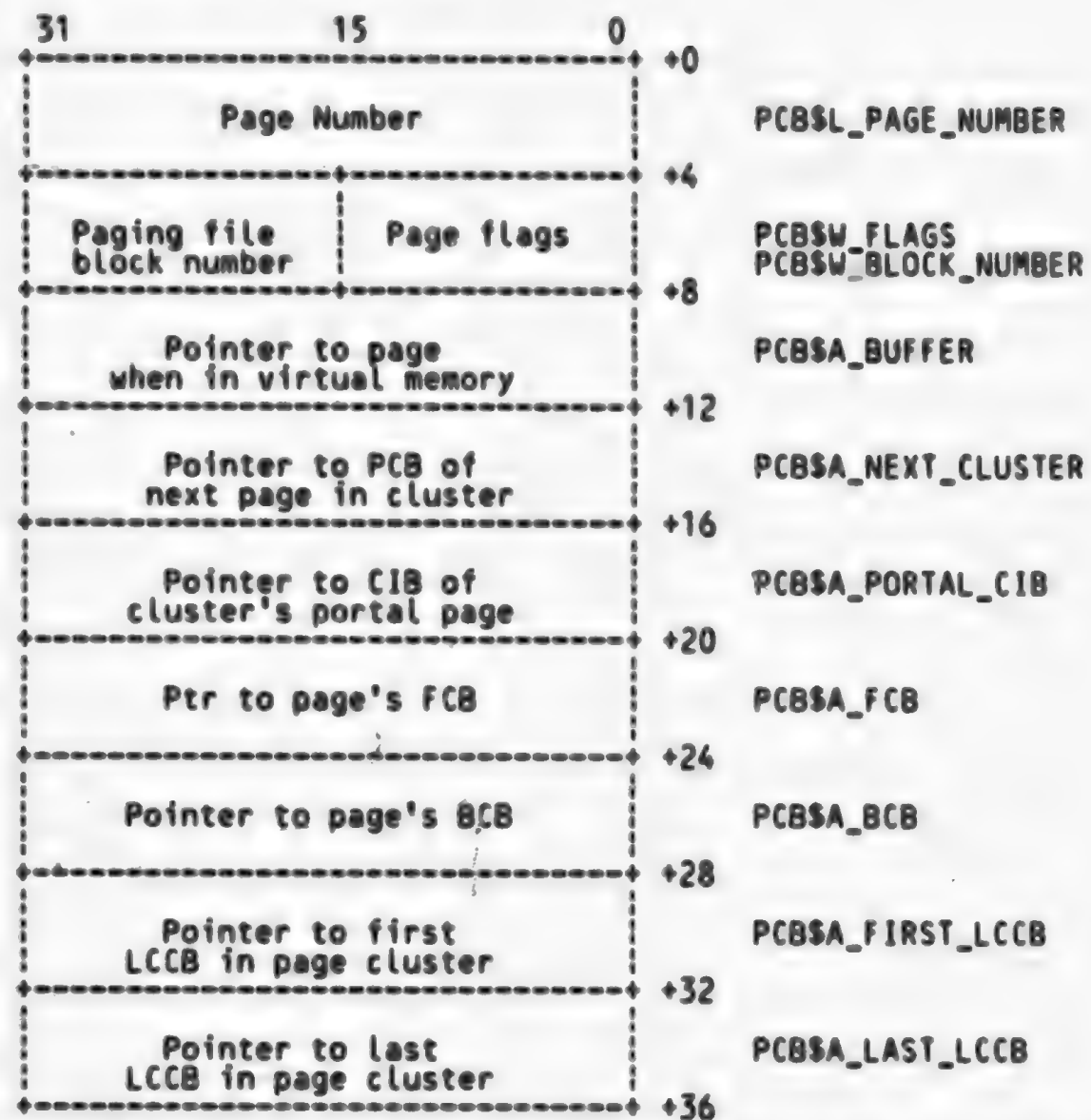
```
CCBSHASH[1, 0, P, S, E] =
```

(CCB\$HASH+CCB\$\$_HASH_BASE+0+(I+HCESS_BLOCK_LENGTH))<P,S,E>;

LITERAL

CCB\$M_CORRUPT = 1^1 - 1^0; ! Stream is corrupt

Page Control Block (PCB)



A Page Control Block (PCB) exists for every page in the cache, and for pages that only have presence locks on them (thus they're not in the staging cache). Portal pages have a Cluster Information Block (CIB) appending to their PCB.

LITERAL

PCBS\$BLOCK_LENGTH = 36;

MACRO

\$PCB = BLOCK(PCBS\$BLOCK_LENGTH, BYTE) FIELD (PCBS\$FIELDS)
\$;

FIELD PCBSZ_FIELDS =
SET

PCBSL_PAGE_NUMBER	=	[0, 0, 32, 0]	
PCBSW_FLAGS	=	[4, 0, 16, 0]	
PCBSV_MODIFIED	=	[4, 0, 1, 0]	Must write back to dict.
PCBSV_NEW_PAGE	=	[4, 3, 1, 0]	Page was in free chain
PCBSV_READ_ONLY	=	[4, 4, 1, 0]	Page is read only
PCBSV_FREE_PAGE	=	[4, 5, 1, 0]	Page is a free page
PCBSV_AVAILABLE	=	[4, 6, 1, 0]	Page available in cache
PCBSV_PORTAL_PAGE	=	[4, 7, 1, 0]	Page is portal page, CIB follows
PCBSW_BLOCK_NUMBER	=	[6, 0, 16, 0]	
PCBSA_BUFFER	=	[8, 0, 32, 0]	
PCBSA_NEXT_CLUSTER	=	[12, 0, 32, 0]	
PCBSA_PORTAL_CIB	=	[16, 0, 32, 0]	
PCBSA_FCB	=	[20, 0, 32, 0]	
PCBSA_BCB	=	[24, 0, 32, 0]	
PCBSA_FIRST_LCCB	=	[28, 0, 32, 0]	
PCBSA_LAST_LCCB	=	[32, 0, 32, 0]	

TES:

LITERAL
PCBSK_LCCB_LIST

= BLOCK[0, PCBSA_FIRST_LCCB; , BYTE];

LITERAL
PCBSM_MODIFIED
PCBSM_NEW_PAGE
PCBSM_READ_ONLY
PCBSM_FREE_PAGE
PCBSM_AVAILABLE
PCBSM_PORTAL_PAGE

= 1^1 - 1^0.	Must write back to dict.
= 1^4 - 1^3.	Page was in free chain
= 1^5 - 1^4.	Page is read only
= 1^6 - 1^5.	Page is a free page
= 1^7 - 1^6.	Page available in cache
= 1^8 - 1^7.	Portal page, CIB follows

Cluster Information Block (CIB)

31	15	0	+0	
	Cluster flags			CIBSW_FLAGS
			+4	CIBSV_REF_COUNTS
Retrieval lock ref count	Presence lock ref count		+8	CIBSW_PRESENCE_COUNT CIBSW_RETRIEVAL_COUNT
Delete lock ref count	Update lock ref count		+12	CIBSW_UPDATE_COUNT CIBSW_DELETE_COUNT
Pointer to first deleted child cluster			+16	CIBSA_FIRST_DELETED_CIB
Pointer to last deleted child cluster			+20	CIBSA_LAST_DELETED_CIB
Ptr to newest lock granted for this page			+24	CIBSA_LAST_LCB
Ptr to oldest lock granted for this page			+28	CIBSA_FIRST_LCB
Ptr to next cluster at this directory level			+32	CIBSA_NEXT_SIBLING
Ptr to prior cluster at this directory level			+36	CIBSA_PRIOR_SIBLING
Ptr to first cluster at next directory level			+40	CIBSA_FIRST_CHILD
Ptr to last cluster at next directory level			+44	CIBSA_LAST_CHILD
Ptr to CIB in next highest directory level			+48	CIBSA_PARENT_CIB
Ptr to portal CIB				CIBSA_HISTORY_CIB

in cluster holding history list	
Ptr to NCB for the node that owns this cluster	CIB\$A_OWNER_NCB
Ptr to cluster's pool header	CIB\$A_POOL

Each clusters' portal page has a CIB associated with it. This block is physically appended to the PCB of the cluster's portal page.

The CIB also points to the cluster's pool header.

The CIB\$V_REF_COUNTS lists must have each of the lock ref counts in the same order as the LOCK\$K_xxx lock request symbols.

LITERAL

CIB\$S_BLOCK_LENGTH = 60 + PCB\$S_BLOCK_LENGTH;

MACRO

%CIB = BLOCK[CIB\$S_BLOCK_LENGTH, BYTE] FIELD (CIB\$Z_FIELDS, PCB\$Z_FIELDS)

%;

FIELD SET CIB\$Z_FIELDS =

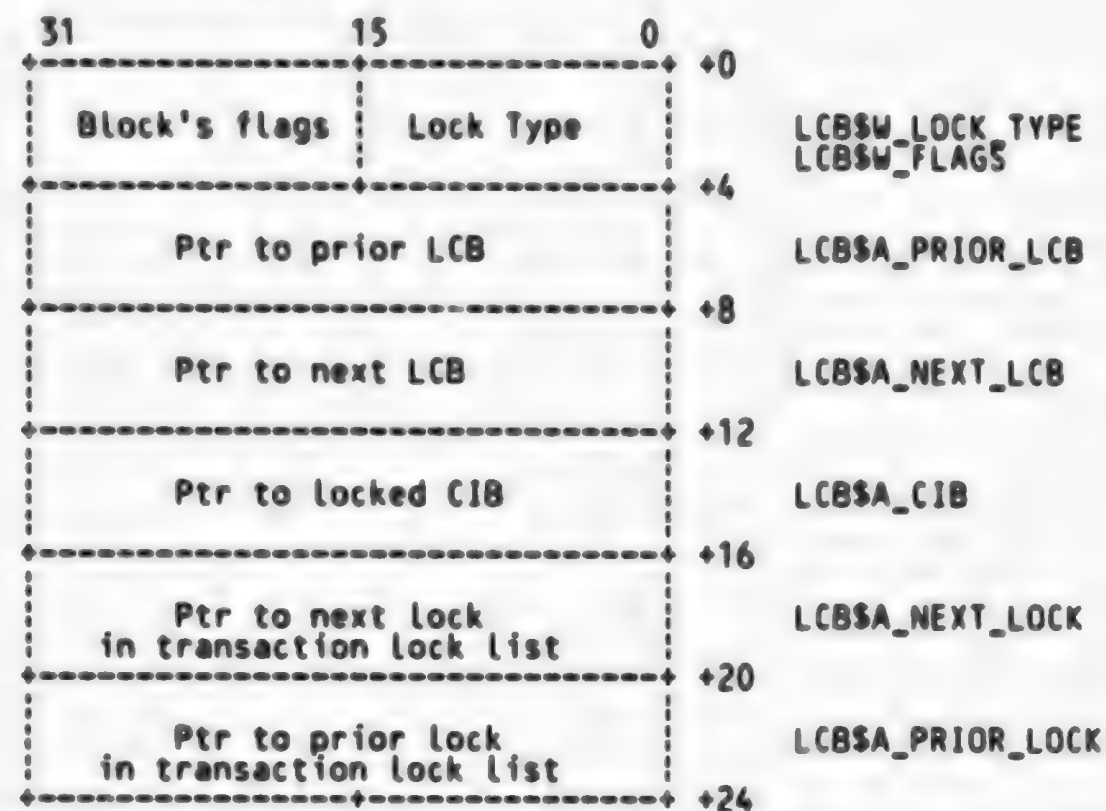
CIB\$W_FLAGS	= [0+PCB\$S_BLOCK_LENGTH, 0, 16, 0],
CIB\$V_LOCKED_DELETE	= [0+PCB\$S_BLOCK_LENGTH, 0, 1, 0],
CIB\$V_LOCKED_UPDATE	= [0+PCB\$S_BLOCK_LENGTH, 1, 1, 0],
CIB\$V_LOCKED_RETRIEVAL	= [0+PCB\$S_BLOCK_LENGTH, 2, 1, 0],
CIB\$V_LOCKED_PRESENCE	= [0+PCB\$S_BLOCK_LENGTH, 3, 1, 0],
CIB\$V_LOCKED	= [0+PCB\$S_BLOCK_LENGTH, 0, 4, 0],
CIB\$V_LOGIN	= [0+PCB\$S_BLOCK_LENGTH, 4, 1, 0],
CIB\$V_COMPLETE	= [0+PCB\$S_BLOCK_LENGTH, 5, 1, 0],
CIB\$V_NEW	= [0+PCB\$S_BLOCK_LENGTH, 6, 1, 0],
CIB\$V_HISTORY	= [0+PCB\$S_BLOCK_LENGTH, 7, 1, 0],
CIB\$V_REF_COUNTS	= [4+PCB\$S_BLOCK_LENGTH, 0, 0, 0],
CIB\$W_PRESENCE_COUNT	= [4+PCB\$S_BLOCK_LENGTH, 0, 16, 0],
CIB\$W_RETRIEVAL_COUNT	= [6+PCB\$S_BLOCK_LENGTH, 0, 16, 0],
CIB\$W_UPDATE_COUNT	= [8+PCB\$S_BLOCK_LENGTH, 0, 16, 0],
CIB\$W_DELETE_COUNT	= [10+PCB\$S_BLOCK_LENGTH, 0, 16, 0],
CIB\$A_FIRST_DELETED_CIB	= [12+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_LAST_DELETED_CIB	= [16+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_LAST_LCB	= [20+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_FIRST_LCB	= [24+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_NEXT_SIBLING	= [28+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_PRIOR_SIBLING	= [32+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_FIRST_CHILD	= [36+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_LAST_CHILD	= [40+PCB\$S_BLOCK_LENGTH, 0, 32, 0],
CIB\$A_PARENT_CIB	= [44+PCB\$S_BLOCK_LENGTH, 0, 32, 0],

```
CIBSA_HISTORY_CIB      = [48+PCBSS_BLOCK_LENGTH, 0, 32, 0];
CIBSA_OWNER_NCB        = [52+PCBSS_BLOCK_LENGTH, 0, 32, 0];
CIBSA_POOL             = [56+PCBSS_BLOCK_LENGTH, 0, 32, 0];
TES;
```

```
LITERAL
CIBSK_LOCK_LIST        = BLOCK[0, CIBSA_LAST_LCB; , BYTE];
CIBSK_SIBLING_LIST     = BLOCK[0, CIBSA_NEXT_SIBLING; , BYTE];
CIBSK_CHILD_LIST       = BLOCK[0, CIBSA_FIRST_CHILD; , BYTE];
```

```
LITERAL
CIBSM_LOCKED_DELETE    = 1^1 - 1^0;
CIBSM_LOCKED_UPDATE    = 1^2 - 1^1;
CIBSM_LOCKED_RETRIEVAL = 1^3 - 1^2;
CIBSM_LOCKED_PRESENCE  = 1^4 - 1^3;
CIBSM_LOCKED           = 1^4 - 1^0;
CIBSM_LOGIN            = 1^5 - 1^4;
CIBSM_COMPLETE         = 1^6 - 1^5;
CIBSM_NEW              = 1^7 - 1^6;
CIBSM_HISTORY          = 1^8 - 1^7;
! Cluster is on login path
! Cluster is complete in cache
! Cluster read in this transaction
! History list is in cluster
```

Lock Control Block (LCB)



LCBs are used to keep track of which locks exist on a cluster (CIB).

Each LCB is linked to its CIB, and to other LCBs for that cluster.
When an LCB is granted, it is placed in the transaction's
lock list until the transaction terminates.

LITERAL

LCBSS_BLOCK_LENGTH = 24;

MACRO

\$LCB = BLOCK[LCBSS_BLOCK_LENGTH, BYTE] FIELD (LCBSZ_FIELDS)
X;

FIELD LCB\$Z_FIELDS =

SET

LCBSW_LOCK_TYPE	= [0, 0, 16, 0];	
LCBSW_FLAGS	= [2, 0, 16, 0];	
LCBSV_CURRENT	= [2, 1, 1, 0];	! Allocated in current transaction
LCBSA_PRIOR_LCB	= [4, 0, 32, 0];	
LCBSA_NEXT_LCB	= [8, 0, 32, 0];	
LCBSA_CIB	= [12, 0, 32, 0];	
LCBSA_NEXT_LOCK	= [16, 0, 32, 0];	
LCBSA_PRIOR_LOCK	= [20, 0, 32, 0];	

TES;

LITERAL
LCBSK_LOCK_LIST
LCBSK_TEMP_LOCK_LIST

= BLOCK(0, LCBSA_PRIOR_LCB; , BYTE),
= BLOCK(0, LCBSA_NEXT_LOCK; , BYTE);

LITERAL
LCBSM_CURRENT

= 1^2 - 1^1; ! Allocated in current transaction

Location Code Control Block (LCCB)



Location codes are used to provide a convenient means for the user to identify a particular object. Each location code is associated with an LCCB. The following objects may be assigned location code:

- 1) Entities
- 2) Lists
- 3) Nodes

LITERAL
LCCBS_BLOCK_LENGTH = 31;

MACRO

\$LCCB = BLOCK[LCCBS\$_BLOCK_LENGTH,BYTE] FIELD (LCCBS\$_FIELDS)
%;

FIELD LCCBS\$_FIELDS =

SET

LCCBS\$_LOCATION_CODE	=	[0. 0. 32. 0].
LCCBS\$_FLAGS	=	[4. 0. 16. 0].
LCCBS\$_AVAILABLE	=	[4. 0. 1. 0].
LCCBS\$_MUST_SCAN	=	[4. 1. 1. 0].
LCCBS\$_GHOST	=	[4. 2. 1. 0].
LCCBS\$_DIRECTORY	=	[4. 3. 1. 0].
LCCBS\$_TERMINAL	=	[4. 4. 1. 0].
LCCBS\$_ENTITY_ATT	=	[4. 5. 1. 0].
LCCBS\$_ENTITY_LIST_ATT	=	[4. 6. 1. 0].
LCCBS\$_ENTITY_LIST	=	[4. 7. 1. 0].
LCCBS\$_STRING_LIST_ATT	=	[4. 8. 1. 0].
LCCBS\$_TYPE	=	[6. 0. 8. 0].
LCCBS\$_LCCB_TYPE	=	[7. 0. 8. 0].
LCCBS\$_OBJECT	=	[8. 0. 32. 0].
LCCBS\$_OBJECT	=	[8. 0. 24. 0].
LCCBS\$_OBJECT	=	[12. 0. 32. 0].
LCCBS\$_PCB	=	[16. 0. 32. 0].
LCCBS\$_NEXT_LCCB	=	[20. 0. 32. 0].
LCCBS\$_PRIOR_LCCB	=	[24. 0. 32. 0].
LCCBS\$_LAST_BLOCK	=	[28. 0. 24. 0].

TES;

LITERAL

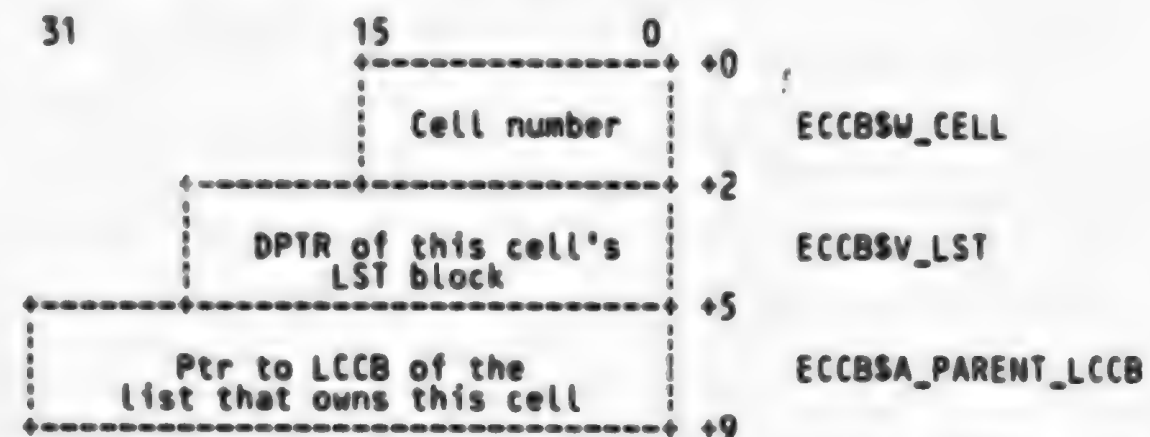
LCCBSK_LCCB_LIST = BLOCK[0, LCCBS\$_NEXT_LCCB; , BYTE];

LITERAL

LCCBSM_AVAILABLE	=	1^1 - 1^0.	! Object's virtual addr is known
LCCBSM_MUST_SCAN	=	1^2 - 1^1.	! Protection tree must be scanned
LCCBSM_GHOST	=	1^3 - 1^2.	! LCCB may not be fetched by LCC
LCCBSM_DIRECTORY	=	1^4 - 1^3.	! Directory NCB
LCCBSM_TERMINAL	=	1^5 - 1^4.	! Terminal NCB
LCCBSM_ENTITY_ATT	=	1^6 - 1^5.	! Entity attribute LCCB
LCCBSM_ENTITY_LIST_ATT	=	1^7 - 1^6.	! Entity list attribute LCCB
LCCBSM_ENTITY_LIST	=	1^8 - 1^7.	! Entity list ECCB
LCCBSM_STRING_LIST_ATT	=	1^9 - 1^8.	! String list LCCB

LCCBSK_LCCB_TYPE_FIRST	=	1.	
LCCBSK_NCB	=	1.	! LCCB includes NCB
LCCBSK_ECCB	=	2.	! LCCB includes ECCB
LCCBSK_LCCB	=	3.	! LCCB
LCCBSK_LCCB_TYPE_LAST	=	3.	

Entity Cell Control Block (ECCB)



Each cell in an entity list may be assigned a location code.

This block is appended to the location code's LCCB to name the specific cell represented by the location code.

LITERAL

```
ECCBS_BLOCK_LENGTH = LCCBS_BLOCK_LENGTH + 9;
```

MACRO

```
%ECCB = BLOCK[ECCBS_BLOCK_LENGTH, BYTE] FIELD (ECCBSZ_FIELDS,  
LCCBSZ_FIELDS)
```

```
%;
```

FIELD ECCBSZ_FIELDS =

```
SET  
ECCBSW_CELL = [0+LCCBS_BLOCK_LENGTH, 0, 16, 0],  
ECCBSV_LST = [2+LCCBS_BLOCK_LENGTH, 0, 24, 0],  
ECCBSA_PARENT_LCCB = [5+LCCBS_BLOCK_LENGTH, 0, 32, 0]
```

```
TES;
```

Node Control Block (NCB)

31	15	0	+0	
User's access rights to this node				NCBSL_ACCESS_RIGHTS
			+4	
Access rights granted by this node's ACL				NCBSL_ACCESS_GRANTED
			+8	
Access rights denied by this node's ACL				NCBSL_ACCESS_DENIED
			+12	
Access rights banished by this node's ACL				NCBSL_ACCESS_BANISHED
			+16	
Pointer to parent NCB block				NCBSA_PARENT_NCB
			+20	
Pswrd's class	Pswrd's data type	Entity's password length		NCBSV_PASSWORD NCBSW_PSW_LENGTH NCBSB_PSW_DTYPE NCBSB_PSW_CLASS
			+24	
Pointer to entity's password string				NCBSA_PSW_POINTER
			+28	
Name's class	Name's data type	Entity name's length		NCBSV_NAME NCBSW_NAME_LENGTH NCBSB_NAME_DTYPE NCBSB_NAME_CLASS
			+32	
Pointer to entity's name string				NCBSA_NAME_POINTER
			+36	
Pointer to Security Preservation Block				NCBSA_SPB
			+40	

The LCCB for a node is followed by an NCB. This gives additional information needed for the node.

LITERAL

NCBSS_BLOCK_LENGTH = 40+LCCBSS_BLOCK_LENGTH;

MACRO

\$NCB = BLOCK(NCB\$\$_BLOCK_LENGTH,BYTE) FIELD (LCCB\$\$_FIELDS,
NCB\$\$_FIELDS)

%;

FIELD NCB\$\$_FIELDS =
SET

NCB\$\$_ACCESS_RIGHTS	=	[0+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_ACCESS_GRANTED	=	[4+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_ACCESS_DENIED	=	[8+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_ACCESS_BANISHED	=	[12+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_PARENT_NCB	=	[16+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_PASSWORD	=	[20+LCCB\$\$_BLOCK_LENGTH, 0, 0, 0]
NCB\$\$_PSW_LENGTH	=	[20+LCCB\$\$_BLOCK_LENGTH, 0, 16, 0]
NCB\$\$_PSW_DTYPE	=	[22+LCCB\$\$_BLOCK_LENGTH, 0, 8, 0]
NCB\$\$_PSW_CLASS	=	[23+LCCB\$\$_BLOCK_LENGTH, 0, 8, 0]
NCB\$\$_PSW_POINTER	=	[24+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_NAME	=	[28+LCCB\$\$_BLOCK_LENGTH, 0, 0, 0]
NCB\$\$_NAME_LENGTH	=	[28+LCCB\$\$_BLOCK_LENGTH, 0, 16, 0]
NCB\$\$_NAME_DTYPE	=	[30+LCCB\$\$_BLOCK_LENGTH, 0, 8, 0]
NCB\$\$_NAME_CLASS	=	[31+LCCB\$\$_BLOCK_LENGTH, 0, 8, 0]
NCB\$\$_NAME_POINTER	=	[32+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]
NCB\$\$_SPB	=	[36+LCCB\$\$_BLOCK_LENGTH, 0, 32, 0]

TES;

File Control Block (FCB)

31	15	0	+0	
File's flags		Channel number		FCBSW_CHAN FCBSW_FLAGS
File Number	Disk pointer to CDDSTOP entity		+4	
			+8	FCBSV_CDDSTOP FCBSB_FILE_NUMBER
First 2 words of the file's FID				FCBSV_FID
			+12	
File ID from lock manager	Last word of file's FID		+16	FCBSW_LOCK_ID
Pointer to first pre-allocated free page				FCBSA_FREE_PAGE
			+20	FCBSV_FILENAME FCBSW_FILE_LENGTH FCBSB_FILE_DTYPE FCBSB_FILE_CLASS
Name's class	Name's data type	Entity's file name length	+24	
Pointer to entity's filename string				FCBSA_FILE_POINTER
			+28	
DKEY of block which points to this file				FCBSL_OWNER
			+32	FCBSV_FULLNAME FCBSW_FULL_LENGTH FCBSB_FULL_DTYPE FCBSB_FULL_CLASS
Name's class	Name's data type	Entity's file name length	+36	
Pointer to entity's full filename string				FCBSA_FULL_POINTER
			+40	
Pointer to temporary free page list				FCBSA_TEMP_FREE
			+44	

Each active (open) dictionary file has an FCB.

The FCBSL_OWNER field holds the DKEY of the attribute block (FIL) which pointed to the file. The primary dictionary file has key 0, while files that are not currently pointed to have a value of -1

! in the owner field.

LITERAL
 FCBSK_FILE_LIMIT = 255;
 FCBSB_BLOCK_LENGTH = 44;

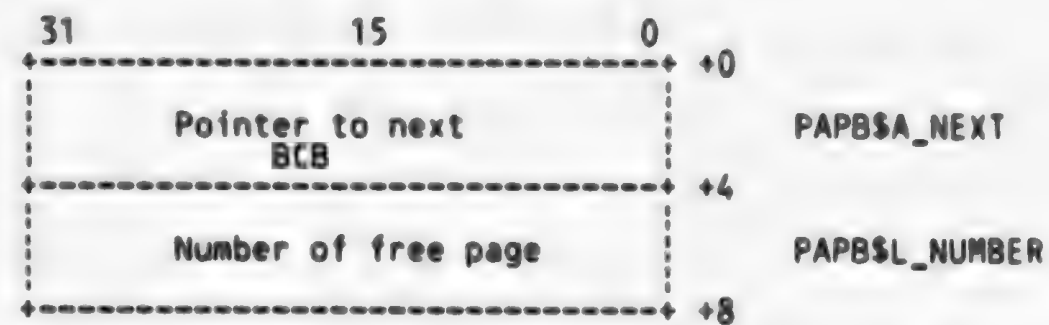
MACRO
 \$FCB = BLOCK[FCBSB_BLOCK_LENGTH,BYTE] FIELD (FCBSZ_FIELDS)
 %;

FIELD SET	FCBSZ_FIELDS =	
	FCBSW_CHAN	= [0. 0. 16. 0].
	FCBSW_FLAGS	= [2. 0. 16. 0].
	FCBSV_READ_ONLY	= [2. 0. 1. 0].
	FCBSV_ROOT	= [2. 0. 1. 0].
	FCBSV_CDDSTOP	= [4. 0. 24. 0].
	FCBSB_FILE_NUMBER	= [7. 0. 8. 0].
	FCBSV_FID	= [8. 0. 0. 0].
	FCBSW_LOCK_ID	= [14. 0. 16. 0].
	FCBSA_FREE_PAGE	= [16. 0. 32. 0].
	FCBSV_FILENAME	= [20. 0. 0. 0].
	FCBSW_FILE_LENGTH	= [20. 0. 16. 0].
	FCBSB_FILE_DTYPE	= [22. 0. 8. 0].
	FCBSB_FILE_CLASS	= [23. 0. 8. 0].
	FCBSA_FILE_POINTER	= [24. 0. 32. 0].
	FCBSL_OWNER	= [28. 0. 32. 0].
	FCBSV_FULLNAME	= [32. 0. 0. 0].
	FCBSW_FULL_LENGTH	= [32. 0. 16. 0].
	FCBSB_FULL_DTYPE	= [34. 0. 8. 0].
	FCBSB_FULL_CLASS	= [35. 0. 8. 0].
	FCBSA_FULL_POINTER	= [36. 0. 32. 0].
	FCBSA_TEMP_FREE	= [40. 0. 32. 0].

TES;

LITERAL
 FCBSM_READ_ONLY = 1^1 - 1^0. ! File can only be opened for read
 FCBSM_ROOT = 1^2 - 1^1. ! FCB is the root dictionary file

Pre-Allocated Page Block (PAPB)



Pre-allocated pages are represented by PAPBs linked onto the FCB in which the pages reside. Each PAPB has the page number of the free page it represents.

LITERAL

PAPBS_BLOCK_LENGTH = 8;

MACRO

\$PAPB = BLOCK[PAPBS_BLOCK_LENGTH, BYTE] FIELD (PAPBSZ_FIELDS)
%:

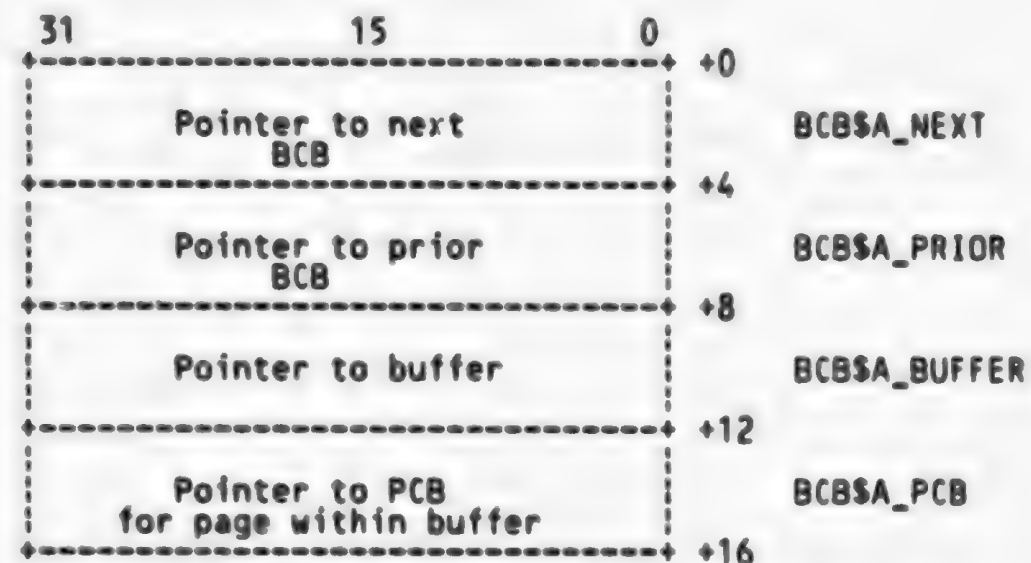
FIELD PAPBSZ_FIELDS =

SET

PAPBSA_NEXT = [0, 0, 32, 0],
PAPBSL_NUMBER = [4, 0, 32, 0]

TES;

Buffer Control Block (BCB)



The buffer control blocks (BCBs) are used to control the pages that are in memory.

The BCBs are linked into a queue. When a buffer is needed, the last buffer in the queue is assigned.

The associated PCB contains information about the page and the buffer. If a page is modified while in the buffer, the page is written back to the work file before the buffer is reused.

LITERAL

```
BCB$K_NUMBER      = 16;    ! Number of in-core buffers
BCB$S_BLOCK_LENGTH = 16;
```

MACRO

```
$BCB = BLOCK[BCB$S_BLOCK_LENGTH,BYTE] FIELD (BCB$Z_FIELDS)
%;
```

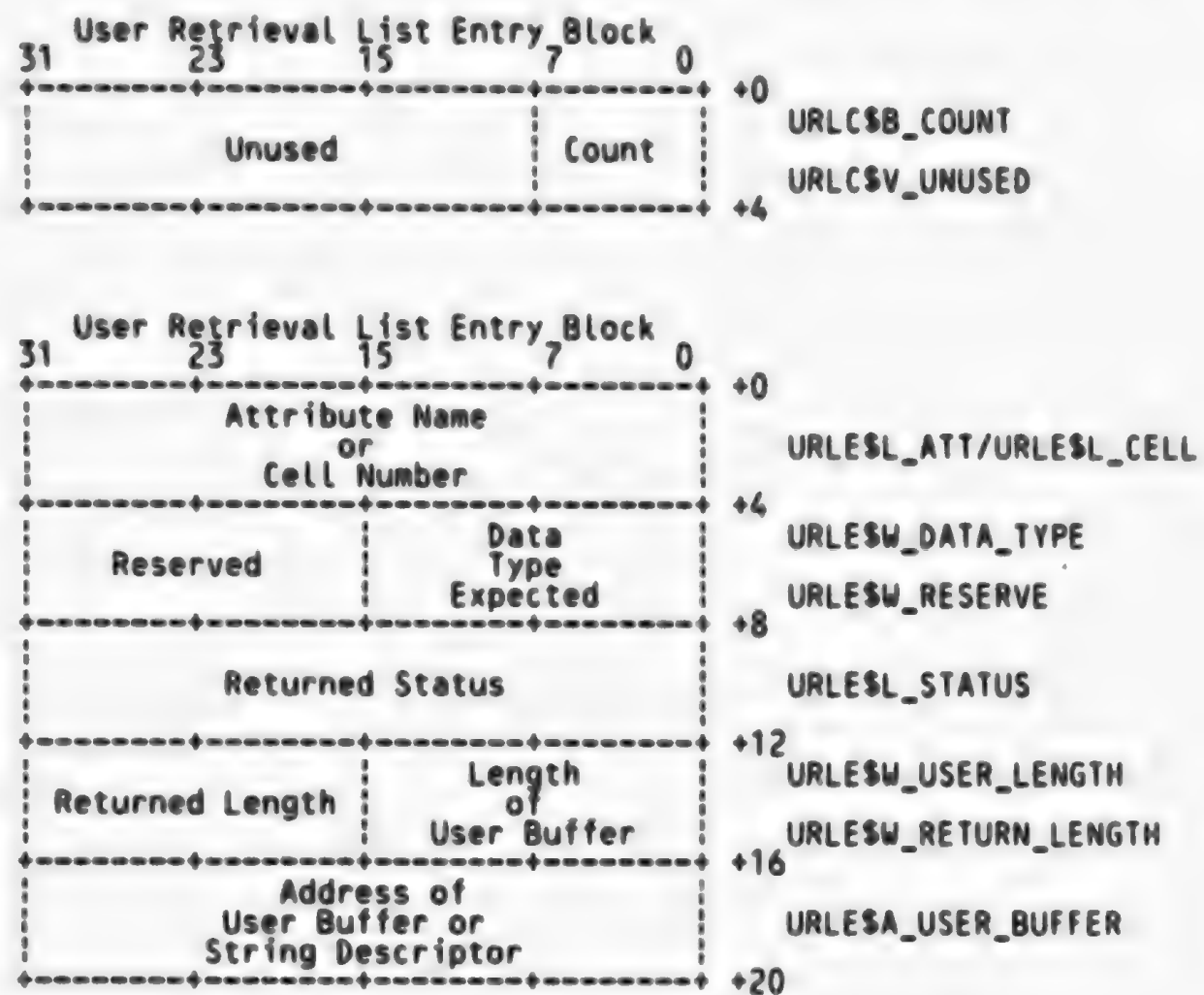
FIELD BCB\$Z_FIELDS =

```
SET
  BCB$A_NEXT      = [0, 0, 32, 0],
  BCB$A_PRIOR     = [4, 0, 32, 0],
  BCB$A_BUFFER    = [8, 0, 32, 0],
  BCB$A_PCB       = [12, 0, 32, 0]
```

```
TES;
```

LITERAL

```
BCB$K_BUFFER_LIST = BLOCK[0, BCB$A_NEXT; , BYTE];
```



LITERAL

```
URLC$S_BLOCK_LENGTH = 4;
URLE$S_BLOCK_LENGTH = 20;
```

MACRO

```
$URLC = BLOCK[URLC$S_BLOCK_LENGTH,BYTE] FIELD (URLC$Z_FIELDS)
%;
```

FIELD URLE\$Z_FIELDS =

```
SET
  URLC$B_COUNT = [0, 0, 8, 0]
  URLC$V_UNUSED = [1, 0, 24, 0]
TES;
```

MACRO

```
$URLE = BLOCK[URLE$S_BLOCK_LENGTH,BYTE] FIELD (URLE$Z_FIELDS)
%;
```

FIELD URLE\$Z_FIELDS =

```
SET
```

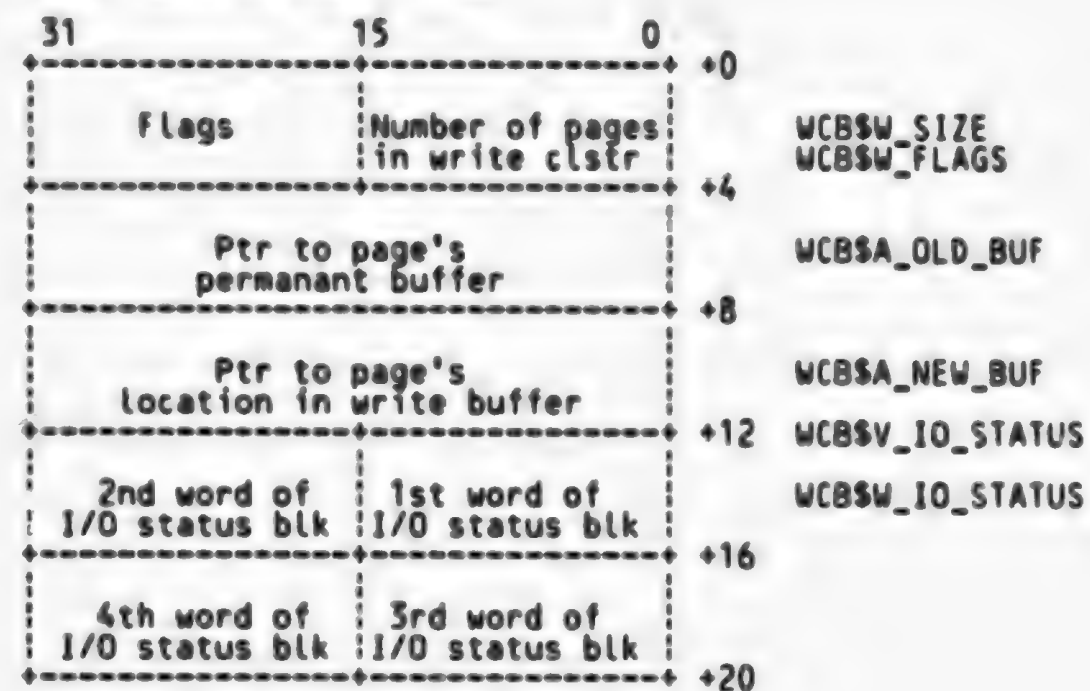
```
URLESL_ATT      = [0, 0, 32, 0],
URLESL_CELL     = [0, 0, 32, 0],
URLE$W_DATA_TYPE = [4, 0, 16, 0],
URLE$W_RESERVE  = [6, 0, 16, 0],
URLESL_STATUS   = [8, 0, 32, 0],
URLE$W_USER_LENGTH = [12, 0, 16, 0],
URLE$W_RETURN_LENGTH = [14, 0, 16, 0],
URLE$A_USER_BUFFER = [16, 0, 32, 0]
```

TES;

STRUCTURE

```
URL$BLOCK [I] = (URL$BLOCK + URLC$S_BLOCK_LENGTH +
(I * URLE$S_BLOCK_LENGTH));
```

Write Control Block (WCB)



LITERAL
WCB\$S_BLOCK_LENGTH = 20;

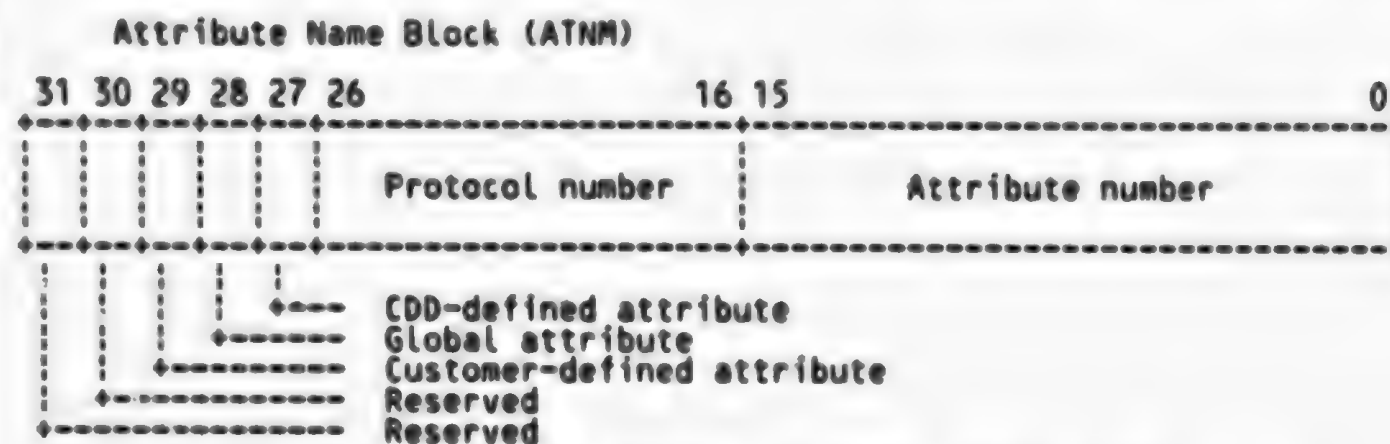
MACRO
\$WCB = BLOCK[WCB\$S_BLOCK_LENGTH, BYTE] FIELD (WCB\$Z_FIELDS)
%;

FIELD SET WCB\$Z_FIELDS =

WCB\$W_SIZE	= [0, 0, 16, 0],	
WCB\$W_FLAGS	= [2, 0, 16, 0],	
WCB\$V_WRITABLE	= [2, 0, 1, 0],	! Write group to file
WCB\$V_GROUP	= [2, 1, 1, 0],	! Group of pages
WCB\$A_OLD_BUF	= [4, 0, 32, 0],	
WCB\$A_NEW_BUF	= [8, 0, 32, 0],	
WCB\$V_IO_STATUS	= [12, 0, 0, 0],	
WCB\$W_IO_STATUS	= [12, 0, 16, 0],	

YES;

LITERAL
WCB\$M_WRITABLE = 1^1 - 1^0, ! Write group to file
WCB\$M_GROUP = 1^2 - 1^1; ! Group of pages



The attribute name block defines the break down of the attribute name. Bits 0 to 15 contains the number. Bits 16 to 26 contains the protocol. The remaining bits are flags define the type of attribute. Bit 27 on implies a system defined attribute. Bit 28 on implies a global defined attribute. Bit 29 on implies a customer defined attribute.

NOTE: If the high order word of the attribute name block is 0 the block is describing a cell. That is the protocol will equal 0 and the flag bits will equal 0.

LITERAL

ATNMSS_BLOCK_LENGTH = 4;

MACRO

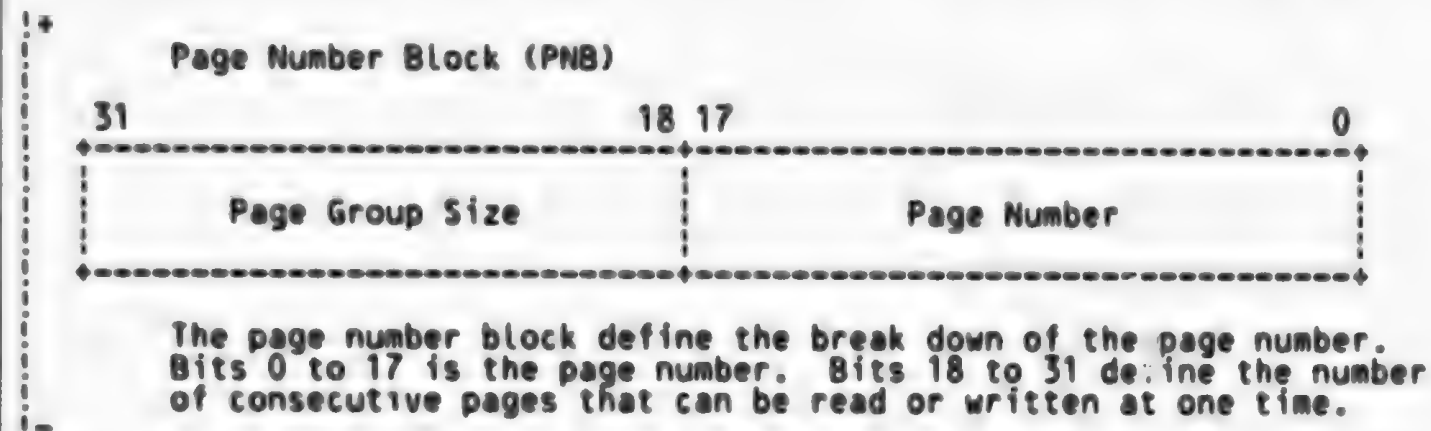
\$ATNM = BLOCK[ATNMSS_BLOCK_LENGTH,BYTE] FIELD (ATNMSZ_FIELDS)
%;

FIELD ATNMSZ_FIELDS =

SET

ATNMSW_NUMBER = [0, 0, 16, 0]
 ATNMSW_HIGH_ORD = [0, 16, 16, 0]
 ATNMSV_PROTOCOL = [0, 16, 11, 0]
 ATNMSV_FLAGS = [0, 27, 5, 0]
 ATNMSV_SYSTEM = [0, 27, 1, 0]
 ATNMSV_GLOBAL = [0, 28, 1, 0]
 ATNMSV_CUSTOMER = [0, 29, 1, 0]

TES;



LITERAL

PNBSS_BLOCK_LENGTH = 4;

MACRO

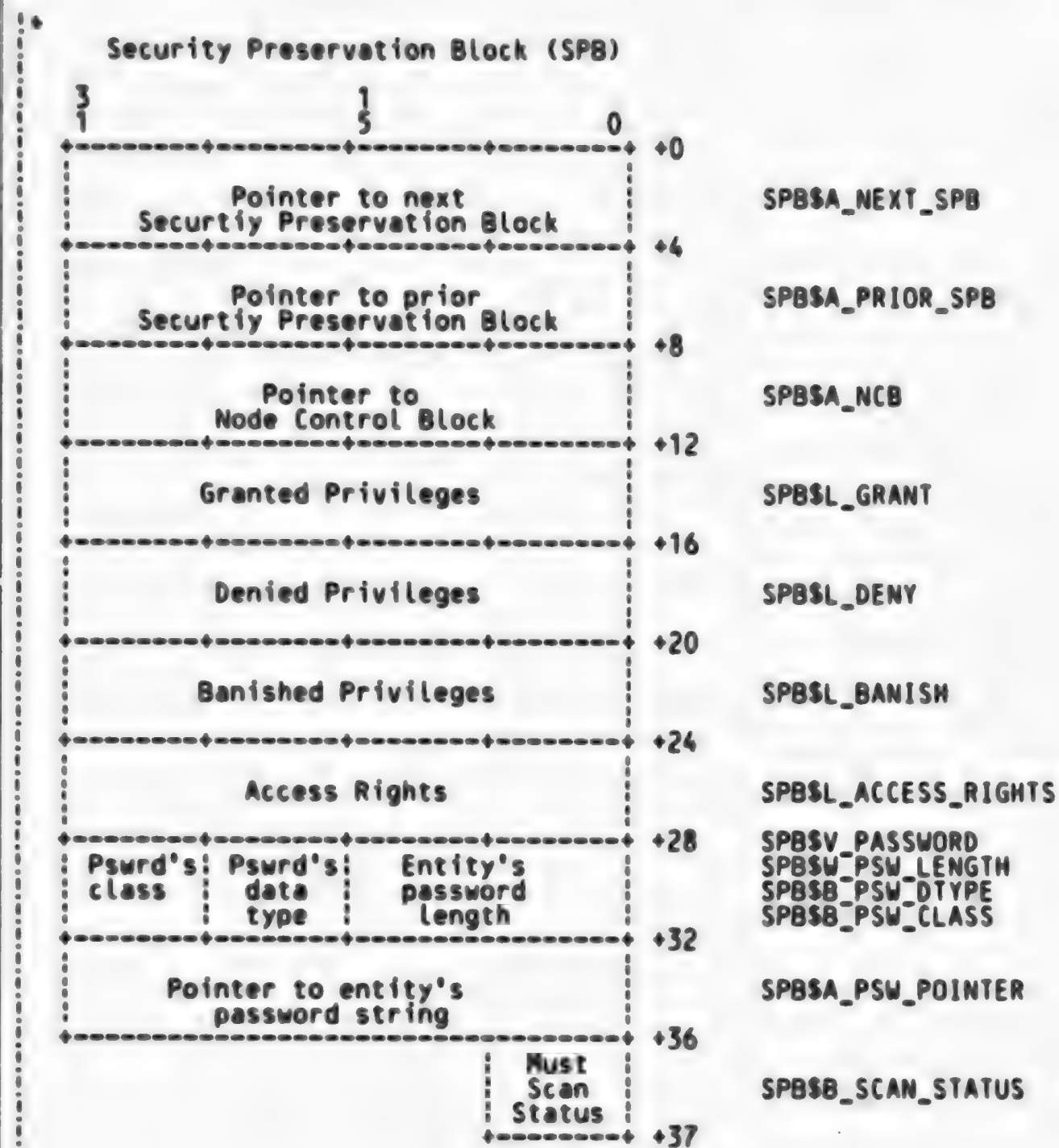
```
$PNB = BLOCK[PNBSS_BLOCK_LENGTH, BYTE] FIELD (PNBSZ_FIELDS)
%;
```

FIELD PNB\$Z_FIELDS =

SET

```
PNB$V_PAGE_NUM = [0, 0, 17, 0],
PNB$V_GROUP_SIZ = [0, 17, 15, 0]
```

TES;



LITERAL
SPB\$S_BLOCK_LENGTH = 37;

MACRO

\$SPB = BLOCK[SPB\$S_BLOCK_LENGTH, BYTE] FIELD (SPB\$Z_FIELDS)
%;

FIELD SPB\$Z_FIELDS =
SET

SPB\$A_NEXT_SPB	=	[0, 0, 32, 0].
SPB\$A_PRIOR_SPB	=	[4, 0, 32, 0].
SPB\$A_NCB	=	[8, 0, 32, 0].
SPB\$S_GRANT	=	[12, 0, 32, 0].
SPB\$S_DENY	=	[16, 0, 32, 0].
SPB\$S_BANISH	=	[20, 0, 32, 0].
SPB\$S_ACCESS_RIGHTS	=	[24, 0, 32, 0].
SPB\$V_PASSWORD	=	[28, 0, 0, 0].
SPB\$W_PSW_LENGTH	=	[28, 0, 16, 0].
SPB\$B_PSW_DTYPE	=	[30, 0, 8, 0].
SPB\$B_PSW_CLASS	=	[31, 0, 8, 0].
SPB\$A_PSW_POINTER	=	[32, 0, 32, 0].
SPB\$B_SCAN_STATUS	=	[36, 0, 8, 0].

TES;

LITERAL

SPB\$K_QUE_HEADER

= BLOCK[0, SPB\$A_NEXT_SPB; , BYTE];

%%SBTTL 'System Literal Definitions'

SYSTEM LITERAL DEFINITIONS

These literals are only used internally.

CDD Implementation Version

LITERAL

CDD\$K_FACILITY	= 43,	
CDD\$K_LOWEST_VERSION	= 201,	! Lowest compatible version
CDD\$K_VERSION	= 201;	! Present version

Boolean literals

LITERAL

TRUE	= 1,	! Boolean TRUE value
FALSE	= 0;	! Boolean FALSE value

The following literals are used to validate routines' parameter lists.

LITERAL

ARG\$K_OPTIONAL	= 1,	! Parameter is optional
ARG\$K_REQUIRED	= 2,	! Parameter is required
ARG\$K_SYNC	= 3,	! See CDD\$SU_VALIDATE documentation
ARG\$K_SYNCIF	= 4,	! See CDD\$SU_VALIDATE documentation
ARG\$K_MARK	= 0,	! Mark parameter as missing
ARG\$K_DEFAULT	= 1,	! Use default value
ARG\$K_STRING	= 1,	! Default value is a null string
ARG\$K_LONG	= 2,	! Default value is a longword
ARG\$K_WORD	= 3,	! Default value is a word value
ARG\$K_REF	= 0,	! Parameter passed by reference
ARG\$K_VALUE	= 1;	! Parameter passed by value

Access Lock Constants

NOTE: The order of the LOCK\$K_NORMAL lock constants MUST be the same as the order of the CIB\$V_REF_COUNTS lock ref count fields in the CIB.

```

LITERAL
LOCKSK_NULL          = 0,          ! Not locked
LOCKSK_NORMAL        = 1,          ! Base of partially queued locks
  LOCKSK_PRESENCE     = 1,
  LOCKSK_RETRIEVAL    = 2,
  LOCKSK_UPDATE       = 3,
  LOCKSK_DELETE       = 4,
LOCKSK_NORMAL_TYPES   = 4,          ! Number of partially queued locks
!!! LOCKSK_QUEUED      = 5,          ! Base of fully queued locks
!!! LOCKSK_FULL_RET    = 5,
!!! LOCKSK_FULL_UPD    = 6,
!!! LOCKSK_QUEUED_TYPES = 2,        ! Number of fully queued locks
LOCKSK_XXX_GET        = 1^17-1^16, ! Flag indicates must establish lock
LOCKSK_XXX_GET_IF     = 1^18-1^17, ! Get lock if not already present
LOCKSK_OPTIONS        = 1^32-1^16, ! Option bits for locks

```

```

!+
!-
Types of page purge requests

```

```

LITERAL
PURGESK_ALL          = 1,          ! Complete purge
PURGESK_PRESENCE     = 2,          ! Purge but keep portal page
PURGESK_RETRIEVAL    = 3,          ! Checkpoint & keep retrieval locks
PURGESK_UPDATE       = 4,          ! Checkpoint & keep all locks
PURGESK_SUBTREE      = 100,        ! Purge whole subtree
PURGESK_CLUSTER      = 101,        ! Purge this cluster only

```

```

!+
!-
Types of blocks that can be allocated in a pool.

```

```

!+
!-
Types of pools

```

```

LITERAL
MEMSK_LOWEST_POOL    = 1,          ! Lowest pool type
  MEMSK_CCB_POOL      = 1,          ! Pool for CCB and HCBs
  MEMSK_CIB_POOL      = 2,          ! Pool for cluster blocks
MEMSK_HIGHEST_POOL   = 2,          ! Highest pool type

```

```

!+
!-
CCB Pool block types

```

```

LITERAL
MEMSK_CCB_LOWEST     = 1,          ! Lowest block type in CCB pool
MEMSK_CCB_CCB        = 1,          ! CCB block

```

```

MEMSK_CCB_LOW_SLOT      = 2,      ! Lowest block type in free block list
MEMSK_CCB_HIGHEST      = 2,      ! HCB block
MEMSK_CCB_HIGHEST      = 2;      ! Highest block type in CCB pool

```

Cluster Pool block types

```

LITERAL
MEMSK_CIB_LOWEST        = 1,      ! Lowest block type in cluster pool
MEMSK_CIB_CIB           = 1,      ! CIB block
MEMSK_CIB_LOW_SLOT     = 2,      ! Lowest block type in free block list
MEMSK_CIB_PCB           = 2,      ! PCB,
MEMSK_CIB_LCB           = 3,      ! LCB,
MEMSK_CIB_LCCB          = 4,      ! LCCB,
MEMSK_CIB_ECCB          = 5,      ! ECCB,
MEMSK_CIB_NCB           = 6,      ! NCB,
MEMSK_CIB_SPB           = 7,      ! SPB,
MEMSK_CIB_HIGHEST      = 7;      ! Highest block type in cluster pool

```

These flags tell the deletion routine how it is to handle the following cases:

```

DELSK_FAST              says that pointers do not have to be cleaned up,
                        as the block they reside in is going to be deleted.

DELSK_PRESERVE          indicates that a directory node is merely to be
                        emptied, and that its cluster is not to be deleted.

DELSK_SUBDICTIONARY     says that sub-files are to have their contents
                        deleted.

```

```

LITERAL
DELSK_FAST              = 1*1 - 1*0,
DELSK_PRESERVE          = 1*2 - 1*1,
DELSK_SUBDICTIONARY    = 1*3 - 1*2;

```

User Identification Criteria

```

LITERAL
CDDSK_ACL_LOWEST       = 1,
CDDSK_ACL_PASSWORD     = 1,      ! PASSWORD
CDDSK_ACL_TERMINAL     = 2,      ! TERMINAL name or class
CDDSK_ACL_UIC          = 3,      ! UIC
CDDSK_ACL_USERNAME     = 4,      ! USERNAME
CDDSK_ACL_HIGHEST      = 4;

```



```

|++  XSBTTL      'Security Masks'
|

```

```

|--- SECURITY MASKS
|

```

```

|+  CDD security bits
|

```

```

| LITERAL
|

```

CDD\$K_PROT_C	= 1^1 - 1^0.	CONTROL access
CDD\$K_PROT_D	= 1^2 - 1^1.	LOCAL DELETE access
CDD\$K_PROT_G	= 1^3 - 1^2.	GLOBAL DELETE access
CDD\$K_PROT_H	= 1^4 - 1^3.	HISTORY list entry creation access
CDD\$K_PROT_P	= 1^5 - 1^4.	PASS THRU access
CDD\$K_PROT_S	= 1^6 - 1^5.	SEE (read) access
CDD\$K_PROT_U	= 1^7 - 1^6.	UPDATE terminal node access
CDD\$K_PROT_X	= 1^8 - 1^7.	EXTEND directory node access
CDD\$K_PROT_F	= 1^9 - 1^8.	FORWARDing directory creation allowed

```

|+  Macro-security values
|

```

CDD\$K_PROT_ANY	= 1^9 - 1^0.
CDD\$K_PROT_DELETE	= CDD\$K_PROT_D OR CDD\$K_PROT_G.
CDD\$K_PROT_EXTEND	= CDD\$K_PROT_F OR CDD\$K_PROT_X.
CDD\$K_PROT_UPDATE	= CDD\$K_PROT_C OR CDD\$K_PROT_D OR CDD\$K_PROT_G OR CDD\$K_PROT_H OR CDD\$K_PROT_U OR CDD\$K_PROT_X OR CDD\$K_PROT_F.

```

|+  Other processor security bits
|

```

```

|+  VAX-11 Datatrieve
|

```

CDD\$K_DTR_PROT_E	= 1^17 - 1^16.	EXTEND file
CDD\$K_DTR_PROT_R	= 1^18 - 1^17.	READ file
CDD\$K_DTR_PROT_M	= 1^19 - 1^18.	MODIFY file
CDD\$K_DTR_PROT_W	= 1^20 - 1^19.	WRITE file

%SBTTL 'User Literal Definitions'

USER LITERAL DEFINITIONS

These symbols are needed by users of the program interface.

System Defined Attribute Names

LITERAL

CDD\$K_SYSNAM_FLAGS = 1^28 OR 1^27 OR 0^16; ! Global/System-defined/Protocol=0

LITERAL

CDD\$K_FIRST_SYSNAM	= 1 OR CDD\$K_SYSNAM_FLAGS,	! Lowest system defined attribute name value
CDD\$K_FILE	= 1 OR CDD\$K_SYSNAM_FLAGS,	! Node's file name
CDD\$K_HISTORY	= 2 OR CDD\$K_SYSNAM_FLAGS,	! History list head
CDD\$K_NAME	= 3 OR CDD\$K_SYSNAM_FLAGS,	! Node's name
CDD\$K_PROTOCOL	= 5 OR CDD\$K_SYSNAM_FLAGS,	! Node's protocol name
CDD\$K_TYPE	= 6 OR CDD\$K_SYSNAM_FLAGS,	! Type of object pointed to by location code
CDD\$K_PATHNAME	= 7 OR CDD\$K_SYSNAM_FLAGS,	! Node's complete pathname
CDD\$K_SHORT_PATHNAME	= 8 OR CDD\$K_SYSNAM_FLAGS,	! Node's path to CDD\$DEFAULT directory
CDD\$K_ORDER	= 9 OR CDD\$K_SYSNAM_FLAGS,	! Directory's order
CDD\$K_LAST_SYSNAM	= 9 OR CDD\$K_SYSNAM_FLAGS;	! Highest system defined attribute name value

Attribute and Entity Types

LITERAL

CDD\$K_FIRST_TYPE	= 1.
CDD\$K_ENTITY	= 1.
CDD\$K_ENTITY_LIST	= 2.
CDD\$K_NULL	= 3.
CDD\$K_NUMERIC	= 4.
CDD\$K_STRING	= 5.
CDD\$K_STRING_LIST	= 6.
CDD\$K_DIRECTORY	= 7.
CDD\$K_TERMINAL	= 8.
CDD\$K_LAST_TYPE	= 8.

User's entity purge options

LITERAL

CDD\$K_ALL	= 1^1 - 1^0.
CDD\$K_ABORT	= 1^2 - 1^1.
CDD\$K_CHECKPOINT	= 1^3 - 1^2.

!+
:- User's node creation options

LITERAL
CDD\$K_NOHISTORY = 1^1 - 1^0, ! Doesn't want history list cluster
CDD\$K_NOACL = 1^2 - 1^1, ! Don't create default ACL entry
CDD\$K_CREATE = 1^3 - 1^2, ! Create dictionary file if needed
CDD\$K_FIRST = 1^4 - 1^3, ! Insert as first node
CDD\$K_LAST = 1^5 - 1^4, ! Insert as last node

!+
:- User's node deletion options

LITERAL
CDD\$K_CHECK = 1^1 - 1^0, ! Fail if directory has children
CDD\$K_SUBDICTIONARY = 1^2 - 1^1, ! Delete contents of subdictionaries

!+
:- Values of the CDD\$K_ORDER attribute

LITERAL
CDD\$K_SORTED = 1, ! Directory is sorted
CDD\$K_NONSORTED = 2, ! Directory is not sorted

%%SBTTL 'LINKAGE DEFINITIONS'

LINKAGE DEFINITIONS

CDDCALL

This linkage uses the CALLG/CALLS linkage convention, except that it allows for one global register to be used in parameter passing.

R11 -- used to pass the user's context pointer.

LINKAGE

CDDCALL = CALL : GLOBAL (USER_CONTEXT = 11);

SYS_JSB

This linkage provides us with a general JSB routine linkage.

LINKAGE

SYS_JSB = JSB;

XSBTTL 'MACRO DEFINITIONS'

MACRO DEFINITIONS

\$ACTIVE
\$INACTIVE

These macros declare that we have started, and finished,
respectively, a CDD transaction. They abort the transaction
if another transaction is in progress.

MACRO
\$ACTIVE =
BEGIN
EXTERNAL
CDD\$GB_INUSE: BYTE;

EXTERNAL LITERAL
CDD\$_NOTASTREE;

BUILTIN
TESTBITSS;

IF TESTBITSS (CDD\$GB_INUSE) THEN
SIGNAL (CDD\$_NOTASTREE);
END
Z,

\$INACTIVE =
BEGIN
EXTERNAL
CDD\$GB_INUSE: BYTE;

CDD\$GB_INUSE = FALSE;
END
Z,

\$BITCLEAR

This macro checks to see if any bit in a mask is set in the
target area. If not, it returns TRUE.

\$BITCLEAR(target, mask) =
(target AND mask) EQLU 0
Z,

\$BITSET

This macro checks to see if any bit in a mask is set in the target area. If so, it returns TRUE.

```
$BITSET(target, mask) =
  (target AND mask) NEQU 0
%,
```

\$DONE_TRANS

This macro is used to terminate a transaction.

Call:

```
$DONE_TRANS [(dsc1 [, dsc2] ...)]
```

Where:

dsci ::= the names of dynamic descriptors which are to have their strings returned to the string pool.

```
$DONE_TRANS (dsc1) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$$$SN_DONE_TRANS      : CDDCALL      NOVALUE;

    CDD$$$SN_DONE_TRANS (dsc1
      %IF NOT %NULL(%REMAINING) %THEN , %REMAINING %FI );
    .USER_CONTEXT[CCBSL_STATUS]
  END
%,
```

\$FIND_ENTITY

This macro returns the virtual address of the LCCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
lccb-block.wa.v = $FIND_ENTITY (valid-arg.ra.v ,
  ( CHECK      ( RETRIEVAL      ) , ( READ      ) );
  ( UPDATE     ( DELETE         ) ) , ( MODIFY   ) );
  DELETE
```

ANY

Where:

valid-arg ::= the address of the calling routine's validated argument list.

The argument list must have the following format:

valid-arg[0] ::= address of longword holding context #

valid-arg[1] ::= address of descriptor holding path name, or zero.

The first set of keywords names the desired lock state of the entity's cluster.

The last set of keywords names the intended access to the cluster.

```

$FIND_ENTITY (valid_arg, locking, security) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$SN_FIND_ENTITY      : CDDCALL;

    CDD$SN_FIND_ENTITY (..valid_arg[1],
      $FIND_XXX_LOCKING (%REMOVE(locking)),
      %NAME('CDD$K_PROT_', security))
  END
%,

$FIND_XXX_LOCKING (class, type) =
  %IF %IDENTICAL (class, %QUOTE LOCK) %THEN
    LOCK$K_XXX_GET OR
  %ELSE
    %IF %IDENTICAL (class, %QUOTE LOCKIF) %THEN
      LOCK$K_XXX_GET_IF OR
    %ELSE
      %IF NOT %IDENTICAL (class, %QUOTE CHECK) %THEN
        %ERROR ('Invalid locking keyword: ', class)
      %FI
    %FI
  %FI
  %NAME ('LOCK$K_', type)
%,

```

\$FIND_NODE

This macro returns the virtual address of the NCB associated with a path name or location code. It also checks to make certain that the cluster is locked as requested, and that the target node allows the requested security access.

Call:

```
ncb-block.wa.v = $FIND NODE (valid-arg.ra.v ,
    ( ( LOCK    RETRIEVAL    READ
      { LOCKIF } { UPDATE    } ) , { MODIFY } ) );
    CHECK    DELETE        DELETE
                                ANY
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.
The argument list must have the following format:

```
valid-arg[0] ::= address of longword holding
                context #
valid-arg[1] ::= address of descriptor holding
                path name, or zero.
valid-arg[2] ::= address of longword holding
                location code, or zero.
```

The first set of keywords names the desired lock state of the node's cluster.

The last set of keywords names the intended access to the cluster.

```
$FIND NODE (valid_arg, locking, security) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$SN_FIND_NODE      : CDDCALL;

    CDD$SN_FIND_NODE (valid_arg,
      $FIND_XXX LOCKING ($REMOVE(locking)),
      $NAME('CDD$K_PROT_', security))
  END
  %.
```

\$FIND_PARENT

This macro returns the virtual address of the NCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
status.wlc.v = $FIND_PARENT (valid-arg.ra.v ,
    ( ( LOCK    RETRIEVAL    READ
      { LOCKIF } { UPDATE    } ) , { MODIFY } ) , ncb-block.wa.r,
    CHECK    DELETE        DELETE
                                ANY
    name.wt.ds);
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.

The argument list must have the following format:

valid-arg[0] ::= address of longword holding context #
 valid-arg[1] ::= address of descriptor holding path name, or zero.
 valid-arg[2] ::= address of longword holding location code, or zero.

The first set of keywords names the desired lock state of the target cluster.

The last set of keywords names the intended access to the cluster.

```
$FIND_PARENT (valid_arg, locking, security, ncb_block, name) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$SSN_FIND_PARENT      : CDDCALL;

    CDD$SSN_FIND_PARENT (valid_arg,
      $FIND_XXX_LOCKING ($REMOVE(locking)),
      $NAME$'CDD$K_PROT_', security), ncb_block, name)
  END
```

z.

\$INIT_DSC

This macro is used to initialize dynamic string descriptors.

Call:

\$INIT_DSC (dsc1 [, dsc2] ...)

```
$INIT_DSC(DSC_NAM) =
  BEGIN
    DSC_NAM[DSC$B_DTYPE] = DSC$K_DTYPE_T;
    DSC_NAM[DSC$B_CLASS] = DSC$K_CLASS_D;
    DSC_NAM[DSC$W_LENGTH] = 0;
    DSC_NAM[DSC$A_POINTER] = 0;
  END
```

z.

\$IO_SYNC (ef, iosb)

This macro waits for the event flag (ef) to be set and for the I/O status parameter (iosb) to be filled in (non-zero).

ef must be the target event flag number.

iosb must be the address of the I/O status block. This must be defined as a VECTOR[WORD] VOLATILE structure.

```

$IO_SYNC (ef, iosb) =
DO
  BEGIN
    LOCAL
      STATUS: LONG;

    STATUS = $WAITFR (efn = ef);
    IF NOT .STATUS THEN
      SIGNAL STOP (.STATUS);
    STATUS = $CLREF (efn = ef);
    IF NOT .STATUS THEN
      SIGNAL STOP (.STATUS);
  END WHILE .iosb[0] EQLU 0
X,

```

\$MARK_PAGE (page-block)

This macro marks a page as modified. Such a page must be written back to the dictionary file when it is purged from the staging buffers.

```

$MARK_PAGE (page block) =
  page_block[PCBSV_MODIFIED] = TRUE
X,

```

\$PARAMETERS (arg1 [, arg2] ...)

This macro is used to build the control vector for the parameter list validate routine (CDD\$SU_VALIDATE).

There is one entry (arg1, arg2, etc) for every formal parameter in the routine. Each entry has the following format:

```

( REQUIRED [, n] [, REF : VALUE] :
  ( SYNC : $SYNCIF ), n :
  OPTIONAL [, MARK : , DEFAULT [, STRING : , LONG : , WORD] ] )

```

```

$PARAMETERS[] =
  UPLIT (%LENGTH, $PARAM_PROCESS (%REMOVE(%REMAINING)))

```

```

%,
SPARM_PROCESS[arg] =
  BYTE ($SPARM_DECIDE ($REMOVE(arg)))
%,
SPARM_DECIDE(status)[] =
  %IF %IDENTICAL (status, %QUOTE REQUIRED) %THEN
    ARGSK_REQUIRED, 0,
    %IF %NULL (%REMAINING) %THEN
      ARGSK_REF, 0
    %ELSE
      $SPARM_REQUIRED (%REMAINING)
    %FI
  %ELSE
    %IF %IDENTICAL (status, %QUOTE SYNC) %THEN
      ARGSK_SYNC, 0, 0, %REMAINING
    %ELSE
      %IF %IDENTICAL (status, %QUOTE SYNCIF) %THEN
        ARGSK_SYNCIF, 0, 0, %REMAINING
      %ELSE
        %IF %IDENTICAL (status, %QUOTE OPTIONAL) %THEN
          ARGSK_OPTIONAL, $SPARM_OPTIONAL (%REMAINING), 0
        %ELSE
          %ERROR ('Invalid parameter status: ', status)
          0, 0, 0, 0
        %FI
      %FI
    %FI
  %FI
%,
SPARM_REQUIRED(number, pass) =
  %IF %NULL (pass) %THEN
    0
  %ELSE
    %NAME ('ARGSK_', pass)
  %FI
  %IF %NULL (number) %THEN
    0
  %ELSE
    , number
  %FI
%,
SPARM_OPTIONAL(action)[] =
  %IF %IDENTICAL (action, %QUOTE MARK) %THEN
    ARGSK_MARK, 0
  %ELSE
    %IF %IDENTICAL (action, %QUOTE DEFAULT) %THEN
      ARGSK_DEFAULT, %NAME ('ARGSK_', %REMAINING)
    %ELSE
      %ERROR ('Invalid action keyword: ', action)
      0, 0
    %FI
  %FI
%,

```

%,

!+
\$PRESENT

This macro checks to see if a parameter is present (non-zero)
in a list.
!-

\$PRESENT (name) =
 .name NEQA 0
%,

!+
\$RECOVERY (RESET ! ENABLE ! DISABLE)

This macro determines the ability of the exit handler to
perform a purge of the cache if the task aborts.

\$RECOVERY (DISABLE)
 declares that the internal data structures or
 external disk structure is in an indeterminant
 state and cannot be recovered by the exit
 handler.

\$RECOVERY (ENABLE)
 declares that the data structure manipulation
 is completed.

\$RECOVERY (RESET)
 specifies that the data structures are in a
 recoverable state.

Note that these can be nested. Recovery is only possible if
the recovery counter is zero (reset).

!-
\$RECOVERY (option) =
 BEGIN
 EXTERNAL
 CDD\$GW_RECOVERY: WORD;
 %IF %IDENTICAL (option, %QUOTE DISABLE) %THEN
 CDD\$GW_RECOVERY = .CDD\$GW_RECOVERY + 1;
 %ELSE
 %IF %IDENTICAL (option, %QUOTE ENABLE) %THEN
 CDD\$GW_RECOVERY = .CDD\$GW_RECOVERY - 1;
 %ELSE
 %IF %IDENTICAL (option, %QUOTE RESET) %THEN
 CDD\$GW_RECOVERY = 0;
 %ELSE
 %ERROR ('Illegal recovery option: ', option)
 %FI
 %FI

%,
END %FI

\$RELEASE_LOCK

This macro calls the CDD\$SN_RELEASE routine to release one or more locks.

Call:

```
$RELEASE_LOCK ( { RETRIEVAL  
                  UPDATE } , ncb-block1 ...);  
                  DELETE
```

\$RELEASE_LOCK (locking)[] =
BEGIN

```
    EXTERNAL ROUTINE  
    CDD$SN_RELEASE      : CDDCALL      NOVALUE;
```

```
    CDD$SN_RELEASE (%NAME ('LOCK$K_', locking), %REMAINING)
```

END
%,

\$SIGNAL_SEVERE (error)

This routine signals a severe error.

\$SIGNAL_SEVERE (ERROR) =
SIGNAL (ERROR OR ST\$K_SEVERE)
%,

\$STATIC_DSC

This macro is used to initialize static string descriptors.

Call:

```
$STATIC_DSC (dsc1 [, dsc2] ...)
```

Where:

```
dsc1 := name ! ( name , source )
```

Source is the name of another descriptor. The named descriptor is initialized to point to the same string as source.


```

$STATIC_DSC[dsci] =
    $STATIC_DSC_BUILD (%REMOVE (dsci))
%,

$STATIC_DSC_BUILD(name, source) =
    BEGIN
        name[DSC$B_DTYPE] = DSC$K_DTYPE_T;
        name[DSC$B_CLASS] = DSC$K_CLASS_S;
        %IF %NULL(source) %THEN
            name[DSC$W_LENGTH] = 0;
            name[DSC$A_POINTER] = 0;
        %ELSE
            name[DSC$W_LENGTH] = .source[DSC$W_LENGTH];
            name[DSC$A_POINTER] = .source[DSC$A_POINTER];
        %FI
    END
%,

```



```

%,
$STRING_PTR_SETUP[PAIR] =
  $STRING_PTR_INIT (%REMOVE(PAIR))
%,
$STRING_PTR_INIT (STR_NAME, STR_VAL) =
  STR_NAME[DSC$A_POINTER] = UPLIT BYTE(%REMOVE(STR_VAL));
%,

```

\$TEXTC

This macro is used to define counted strings. The first byte of such strings is a count of the number of characters in the string.

Each string is defined to be a VECTOR[BYTE] structure, with the 0 element being the character count, and the actual string starting at STRING[1].

```
$TEXTC ((name1,'str1') [, (name2, 'str2')] ...);
```

```

$TEXTC[PAIR] =
  $TEXTC_STR(%REMOVE(PAIR))
%,
$TEXTC_STR(NAME,TSTR) =
  BEGIN
    NAME = UPLIT BYTE (%CHARCOUNT(%REMOVE(TSTR)), %REMOVE(TSTR)) :
    VECTOR[%CHARCOUNT(%REMOVE(TSTR))+1, BYTE];
%,

```

\$VALIDATE (cntl, [arg-list])

This macro generates a call to the general transaction setup routine.

cntl -- the address of the control vector for CDD\$\$U_VALIDATE.

arg-list-- is optional. If present, it is the address of the vector which is to receive the verified argument list from CDD\$\$U_VALIDATE.

```

$VALIDATE(cntl, arg_list) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$$SN_START_TRANS      : CDDCALL NOVALUE;

```

BUILTIN
AP;

%IF %NULL (arg_list) %THEN
CDD\$\$\$SN_START_TRANS (.AP, cntl)
%ELSE
CDD\$\$\$SN_START_TRANS (.AP, cntl, arg_list)
%FI

%; END

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